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FEABL (FINITE ELEMENT ANALYSIS BASIC LIBRARY) USER'S GUIDE

Oscar Orringer, et al

Massachusetts Institute of Technology

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FEABL
(FINITE ELEMENT ANALYSIS
BASIC LIBRARY)
USER'S GUIDE

Oscar Orringer
Susan E. French

AEROELASTIC AND STRUCTURES RESEARCH LABORATORY
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MASSACHUSETTS INSTITUTE OF TECHNOLOGY
CAMBRIDGE, MASSACHUSETTS 02139

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13. ABSTRACT

This guide contains complete instructions for the use of FEABL, a basic software system developed for finite element analysis at the MIT Aeroelastic and Structures Research Laboratory. FEABL has been designed primarily for (but is not limited to) the specialized type of continuum analysis problem encountered in the materials laboratory. FEABL has also been used successfully as an educational tool in the finite element analysis course given by the MIT Department of Aeronautics and Astronautics. The software is modular, and is written in machine-independent FORTRAN IV. A complete program listing is contained in Appendix C. Element subroutines which can be used either independently or in conjunction with FEABL will be presented in future publications.

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FOREWORD

ABSTRACT

This guide contains complete instructions for the use of FEABL, a basic software system developed for finite element analysis at the MIT Aeroelastic and Structures Research Laboratory. FEABL has been designed primarily for (but is not limited to) the specialized type of continuum analysis problem encountered in the materials laboratory. FEABL has also been used successfully as an educational tool in the finite element analysis course given by the MIT Department of Aeronautics and Astronautics. The software is modular, and is written in machine-independent FØRTRAN IV. A complete program listing is contained in Appendix C. Element subroutines which can be used either independently or in conjunction with FEABL will be presented in future publications.

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SECTION 1

INTRODUCTION

1.1   Purpose and Scope

Finite element analysis programs developed in the past
fall into two general categories: the "one-shot" program
optimized for solution of a particular problem, or the "systems"
program designed for the non-optimum solution of any problem.
The structural engineer who has a new analysis to perform must
spend a great amount of time either writing his own "one-shot" pro-
gram or learning how to use one of the general "systems". FEABL
attempts to relieve the analyst of these burdens by providing
a basic library which handles the numerical operations common
to all finite element analysis, and which can be understood
and used after a few hours of study. FEABL's current capa-
bilities are limited to static analysis. Other options will
be added in the future, as the need arises, and according
to their usefulness. This guide is written for the analyst
who has some acquaintance with the theory of finite element
analysis and with IBM FØRTRAN IV.

THE FEABL software system has been designed primarily for
"materials laboratory" stress analysis. The analysis problems
encountered in this application area are characterized by geometry
which, although not simple, is mathematically describable, and
by stress solution accuracy requirements which can usually be met
with models of 1,000 or fewer degrees of freedom. Stress concen-
tration analysis for two-material systems and for composite-to-
metal step-lap joints are two examples of the materials laboratory
type of problem. Another feature of many of these problems is a
requirement for nonstandard solution methods. For instance,
the boundary conditions may include an undetermined prescribed
displacement, the value of which depends upon satisfaction of an

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auxiliary condition on a boundary stress integral; a specific example is given in Ref. 1. The main features of FEABL, all-in-core computation and control of the solution process by the user, are well adapted to the needs of such problems.

Experience in the MIT Department of Aeronautics and Astronautics has also shown that FEABL is a valuable educational tool. A budget of $50 to $70 per man on an IBM 370/155 machine allows each student to solve one or two problems, each having 200 to 300 degrees of freedom. In the approach taken at MIT, the student is required to program his structure geometry, element interconnections, and subroutines for generating stiffnesses and stresses. He must also achieve a general understanding of the principles upon which the FEABL subroutines are based. Thus, each student is able to gain some practical experience, as well as theoretical knowledge in one semester of a finite element analysis course. This approach is felt to be superior to the traditional method of having the student learn only how to code the input data for a particular, complete "systems" program.

The flexibility and economy of the FEABL software system has been amply demonstrated by various applications in which real problems were solved during the test/exercise development phase. The following are examples abstracted from these applications:

Problem Description	Total DOF	CPU Time (Min)	Total Cost ($)
Displacement analysis of two large rigid frame structures (6 load cases each)	216 426	0.42 2.49	5.37 13.17
Continuum stress concentration analysis with boundary condition determined by auxiliary stress integral (3 complete solutions required for each case)	338 442	1.50 2.70	4.50 12.00
Stress analysis of a large frame in an oil tanker, using rectangle and triangle continuum elements and flange elements (6 loading cases)	1,266	3.0 (Approx.)	14.00
Pressure distribution analysis for viscous flow around a corpuscle in a blood vessel	160	No data available	
Stress analysis, including solution for stress intensity factor of a sharp crack, using hybrid element at the crack	160	0.13	2.54

The above examples were run on the IBM 370/155 at the MIT Information Processing Center. The costs presented are for production runs, exclusive of debugging.

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1.2   General System Concepts

Development of the FEABL software has been based on three
general concepts. First, FØRTRAN IV was selected as the program
language, and all FØRTRAN syntax rules were followed rigorously.
This makes FEABL as machine-independent as possible, to allow
for use at installations having other than IBM equipment. The
only restrictions are that the FØRTRAN compiler software must
permit logical IF statements, variable dimensioning, six-character
variable names, and labelled CØMMØN declarations. Second, the
software is modular in character. The finite element analysis
process has been decomposed into a series of distinct steps,
and each step which is independent of structure geometry has been
programmed as a separate subroutine. FEABL is thus analogous,
in a sense, to the IBM Scientific Subroutine Package.

The third concept deals with data storage techniques. In
order to minimize wasted storage space, a vector approach has
been adopted. All of the input data for a problem, certain
"housekeeping" data, and internally generated problem data such
as the master stiffness matrix are stored in a single vector,
referred to below as the /DATA/ vector. In addition, control
parameters which locate entries in the /DATA/ vector or manage
input-output operations are organized in four CØMMØN groups.

1.3   General System Organization

The analyst must provide one or more programs which inter-
face with the FEABL software:
1. A MAIN program which inputs the required problem data
   and controls the execution of all sub programs.
2. One or more subroutines which generate an element
   stiffness matrix (and, if required, equivalent nodal
   forces for thermal stress, gravity, etc.) for each
   type of element.*

---

*
Subroutines from ASRL EGL (Element Generator Library) may be
used.

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3. One or more subroutines which generate element stresses
    from the element's nodal displacements.[*]

Data may be input by reading cards, disk or tape files, by auto-
matic generation techniques, or by a combination of the two
methods. The details of a particular problem will determine
which approach is more efficient. For very simple types of prob-
lems, in which the structure is divided into only a few different
element types and shapes, stiffness matrix and stress generation
may be done within the MAIN program. However, if separate sub-
routines are required for these tasks, pre-existing ones may be
interfaced with FEABL, since the FEABL software package has been
written in a manner such that pre-existing FØRTRAN subroutines
may be adapted to run with FEABL with little or no re-programming.

The interface between user-written programs and FEABL
consists of two parts: data location and process control.
Figure 1 illustrates the general features of the data location
interface. The labelled CØMMØN areas mentioned in the previous
section must appear in the user's MAIN program. This is accom-
plished simply by including identical sets of DIMENSIØN, CØMMØN
and EQUIVALENCE declarations in MAIN and in the FEABL software.
There is no requirement for the element generator subroutines
to communicate with FEABL in this manner; these subroutines need
only be interfaced with user's MAIN by an identical DIMENSIØN
declaration for the element stiffness matrix and "force/displace-
ment" vector. The latter is a dual-purpose vector in which
element equivalent nodal forces are generated, and which serves
later as a storage area for the solution displacement vector of
the element. The interface between the generator routines and
FEABL is achieved automatically by variable DIMENSIØN declarations
contained in the FEABL software. The data location interface is
discussed in detail in Section 2.

---

[*]Subroutines from ASRL EGL (Element Generator Library) may be
used.

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FIGURE 1

The process control interface is simply a recognition
that input of problem data and operations on the data must be
performed in a certain sequence. This requires a definite series
of programming steps in user's MAIN program. Section 3 discusses
the process control interface in detail.

THE FEABL DATA STORAGE SYSTEM

2.1   Accessory Data

The accessory data is divided into four control parameter
CØMMØN groups. The CØMMØN declaration statements listed below
for the four control parameter groups must appear in all user-
written MAIN programs.

2.1.1 Input-Output Control Parameters

CØMMØN /IØ/ KR, KW, KP, KT1, KT2, KT3

The six parameters in the /IØ/ group are hardware device
codes, defined as follows:

KR - card reader

KW - printer

KP - card punch

KT1, KT2, KT3 are extra device codes which the user may define
optionally as he pleases. For example, KT1 might refer to a
system direct access disk, KT2 to an external tape unit, etc.
Only the printer device code KW is used directly by the FEABL
software presented in this guide.   The input-output control
parameters are used in FØRTRAN READ and WRITE statements, e.g.:

```
      READ (KR,1) NET, NDT
   1  FØRMAT (2I6)
      WRITE (KW,2) NET, NDT
   2  FØRMAT (16HOTØTAL ELEMENTS=, I6, 10X, 10HTØTAL DØF=, I6)
```

The correct values of KW and any other device codes employed by
the user must be established at the beginning of the MAIN program.

2.1.2 Problem Size Parameters

CØMMØN /SIZE/ NET, NDT

These parameters define the problem size in terms of the
total number of elements (NET) and the total number of degrees
of freedom (NDT) in the whole structure. NDT includes both

constrained and unconstrained degrees. In standard solutions
these parameters are established at the beginning of the MAIN
program and remain fixed.

2.1.3 Address Index Parameters

```
CØMMØN /BEGIN/ ICØN, IKØUNT, ILNZ, IMASTR, IQ, IK
CØMMØN /END/ LCØN, LKØUNT, LLNZ, LMASTR, LQ, LK
```

These parameters control the begin and end locations of
data sub-blocks in the /DATA/ vector. The user establishes cor-
rect values for these parameters by calling FEABL subroutine
SETUP after he has established the /SIZE/ block. Use of the
address index parameters to locate entries in the /DATA/ vector
is discussed in Subsection 2.2.

2.2 The /DATA/ Vector

```
DIMENSIØN REAL(xxxx), INTGR(xxxx)
CØMMØN /DATA/ REAL
EQUIVALENCE (REAL (1), INTGR(1))
```

The above declarations serve to define the /DATA/ vector
as a one-dimensional array occupying a CØMMØN area labelled
/DATA/ and having two reference names: REAL for floating point
entries and INTGR for integer entries. These declarations must
appear in the user-written MAIN program. The user chooses the
dimension integer "xxxx" to suit his particular needs. The
DIMENSIØN declaration with user's value for the length of the
/DATA/ vector must be inserted in each FEABL subprogram, immedi-
ately following the subroutine name declaration, e.g.:

```
SUBRØUTINE ASEMBL (LNUM, NDE, ELK, ELQ)
DIMENSIØN REAL(1000), INTGR(1000)
```

Some compilers restrict the length of any vector to
32,768 words ($8,000_{HEX}$). If this restriction is encountered,
and the user desires a longer /DATA/ vector, say 50,000 words,
the following form of the declaration statements can be employed:

```
DIMENSIØN REAL(25000), INTGR(25000), DUMMY(25000)
CØMMØN /DATA/ REAL, DUMMY
EQUIVALENCE (REAL(1), INTGR(1))
```

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The CØMMØN /DATA/ declaration must then be changed to the above form in each FEABL subprogram.

The dual nature of the /DATA/ vector must be recognized clearly. For example, the 100th entry in the vector may be treated as either a floating point or an integer quantity in arithmetic or logical instructions by operating, respectively, with REAL(100) or INTGR(100). This property will lead to compilation errors if the user attempts to save storage space by declaring INTGR to be an array of 2-BYTE words (half-words), while REAL is left as an array of normal 4-BYTE words (full words), or if the user attempts to run in double precision mode for floating point arithmetic.

The /DATA/ vector is organized into six blocks, the limits of which are determined by the values stored in the address index parameters (Subsection 2.1.3). In normal usage, the first four blocks contain integer data and the last two contain floating point data. Storage conventions are detailed in the following sections, in the order in which the blocks appear in the /DATA/ vector.

2.2.1 Constraint Vector

This is an integer block which contains the global number of each degree of freedom at which a displacement is to be prescribed. Entries in the constraint vector are referred to by:

INTGR(I) where ICØN$\leq$I$\leq$LCØN

The following example illustrates user action involving the constraint vector

```
    .
    .
    .
INTGR(ICØN)=2
INTGR(ICØN+1)=25
INTGR(LCØN)=6
    .
    .
```

By means of the above instructions the user has specified that global displacements 2, 25, and 6 will be prescribed. The prescribed displacements need not be listed in any particular order. Any excess space in the constraint vector is filled with zeros by FEABL.

2.2.2 Address Count Vector

This is an integer block which contains information relating to the absolute address in the /DATA/ vector of each diagonal entry of the master (global) stiffness matrix. Entries in the address count vector are referred to by:

INTGR(I) where $IK\emptyset UNT \leq I < LK\emptyset UNT$

To obtain the relevant data for the Nth row, the correct subscript is:

$J = IK\emptyset UNT + N - 1$

INTGR(J) contains address information for row N

Normally, the user has no direct communication with the address count vector; its contents are generated internally by FEABL. The absolute address of each diagonal entry is output for use in debugging. After the output, the data in the address count vector are modified as follows:

Contents = (Absolute Address of K_{ii}) - i

Subsequently, any <u>stored</u>[*] entry K_{ij} can be obtained by referring to:

REAL(KADR)

where the address is given by:

$KADR = INTGR(IK\emptyset UNT + I - 1) + J$

2.2.3 Leading Non-Zero Entry Vector

The LNZ vector is an integer block containing the number j of the column in which appears the leading non-zero entry K_{ij}

---

[*]Many entries of the master stiffness matrix are not stored in the /DATA/ vector. See Subsection 2.2.6 for details.

of the master stiffness matrix for each row i. Entries in the
LNZ vector are referred to by:

$$INTGR(I) \text{ where } ILNZ \leq I \leq LLNZ$$

INTGR(ILNZ+N-1) refers to row N.

Normally, the user has no direct communication with the
LNZ vector; its contents are generated internally by FEABL, and
are output for use in debugging and evaluating displacement
numbering strategies.

2.2.4 Master Assembly List

This is an integer block containing all of the information
which relates the user's local degree of freedom numbers to the
user's global numbering system. The master assembly list is, in
effect, a Boolean logical transformation between the element
displacement vectors, which together form a linearly dependent
set, and the global displacement vector, which is an independent
set. Entries in the master assembly list are referred to by:

$$INTGR(I) \text{ where } IMASTR \leq I \leq LMASTR$$

The master assembly list need not be filled, but if it is not,
a zero must be stored immediately following the last active
position.

The master assembly list is subdivided into two sections,
as indicated in Figure 2.

FIGURE 2

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The pointer section consists of one location per element, from INTGR(IMASTR) to INTGR(IMASTR+NET-1), with the convention that these locations correspond to the user's elements in ascending order 1, 2,..., NET. Each pointer contains the absolute address in the /DATA/ vector of the location where the assembly list for an element starts, as indicated in Figure 3.

FIGURE 3

The remainder of the master list contains the user's element-by-element sequence of global displacement numbers. As an illustrative example, consider the set of elements shown in Figure 4, consisting of two 4-node rectangles and two 3-node triangles, each with two degrees of freedom per node. The total length required for the master assembly list is given by:

```
Pointers:    1 per element-----------4
Rectangles:  8 DOF/element----------16
Triangles:   6 DOF/element----------12
TOTAL----------------------------------32
```

Suppose that IMASTR=101. Then the assembly information, correctly stored, would appear as shown in Figure 5.

= Element Global Number

= Global DOF Numbers

Local Numbering
Conventions for DOF

FIGURE 4

FIGURE 5

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The following crude set of FØRTRAN instructions might be used to store the assembly information shown above:

```
NEXT=IMASTR+4
INTGR(IMASTR)=NEXT    (establishes pointer for 1st element)
INTGR(NEXT)=1
INTGR(NEXT+1)=2
   .
   .
   .
INTGR(NEXT+7)=4
NEXT=NEXT+8
INTGR(IMASTR+1)=NEXT (establishes pointer for 2nd element)
INTGR(NEXT)=5
INTGR(NEXT+1)=6
   .
   .
   .
INTGR(NEXT+5)=8
NEXT=NEXT+6
INTGR(IMASTR+2)=NEXT (establishes pointer for 3rd element)
   .
   .
   .
```

The use of pointers in the master list allows assemblies involving as many different types of elements with different total numbers of degrees of freedom as desired. Also the use of DOF numbers rather than node numbers permits assembly of elements having different numbers of displacements at various nodes without requiring any special programming or conventions. The user may choose any element and displacement global numbering schemes and any local displacement numbering conventions he desires. The only restrictions are that:

1. There must be no gaps in the master assembly list.

2. The array must be filled or, if excess storage exists, a zero must be stored after the last displacement number.

3. All element numbers and degree of freedom numbers must be positive.

4. The lowest element number and the lowest global degree of freedom number must each be unity.

2.2.5 Force/Displacement Vector

This is a floating point block which contains all of the
force and displacement information required for analysis of a
structure. Entries in the force/displacement vector are referred
to by:

REAL(I) where IQ$\leq$I$\leq$LQ

The Nth entry in the vector is the force or displacement associ-
ated with the Nth global degree of freedom.

Various types of information are <u>overlayed</u> in the force/
displacement vector. First, if the structure being analyzed is
loaded by continuum body forces (e.g., gravity) or is in a thermal
environment, the resulting <u>element equivalent nodal forces</u> must
be assembled along with the element stiffness matrices. FEABL
subroutine ASEMBL uses the force/displacement vector as a storage
area for this purpose. Second, the user must introduce his glo-
bal concentrated nodal forces and prescribed displacements into
the force/displacement vector. Finally, the FEABL solution sub-
programs store the displacement solution in the force/displac-
ment vector by overwriting the prescribed quantities.

The following simple algorithm enables the user to com-
municate with the Nth global degree of freedom:

        NN=IQ+N-1
        REAL(NN)=...
        or
        ...=f(REAL(NN))

For example, suppose the structure in Figure 4 is to be given
the boundary conditions shown in Figure 6. Global degrees of
freedom 4 and 8 have concentrated forces A and B applied, respec-
tively, while degrees 12 and 15 have displacements C and D
prescribed, respectively. Displacements, 1, 2, 5, 6, 13, and 14
are prescribed to be zero. The following set of FORTRAN
instructions specify these boundary conditions:

(a) Establishment of the constraint vector

```
INTGR(ICØN)=1
INTGR(ICØN+1)=2
INTGR(ICØN+2)=5
INTGR(ICØN+3)=6
INTGR(ICØN+4)=12
INTGR(ICØN+5)=13
INTGR(ICØN+6)=14
INTGR(ICØN+7)=15
```

(b) Input of the prescribed displacements:

```
REAL(IQ)=0.
REAL(IQ+1)=0.
REAL(IQ+4)=0.
  .
  .
  .
REAL(IQ+13)=0.
REAL(IQ+11)=C
REAL(IQ+14)=D
```

(c) Input of the prescribed forces:

```
REAL(IQ+3)=REAL(IQ+3)+A
REAL(IQ+7)=REAL(IQ+7)+B
```

The instructions (b) and (c) have been written assuming that element equivalent nodal forces have been assembled into the force/displacement vector. Thus, all degrees at which displacements are prescribed must be set to their correct values, while nonzero concentrated global forces must be added to the pre-existing assembled element forces.

A, B are global concentrated forces

C, D are global prescribed displacements

FIGURE 6

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2.2.6 Master Stiffness Matrix

This floating point block contains the essential entries of the master stiffness matrix, stored one row after the other. Entries in the master stiffness matrix are referred to by:

REAL(I) where $IK \leq I < LK$

Normally, the user is not required to communicate directly with the master stiffness matrix, but an understanding of its detailed organization will be helpful in debugging.

Since a stiffness matrix is always symmetric, only its lower triangle need be kept in storage. Thus, the general organization of the master stiffness array in the /DATA/ vector may be represented by the diagram shown in Figure 7.

FIGURE 7

Significa.it additional savings in storage may be realized for problems with many degrees of freedom by taking advantage of the fact that a master stiffness matrix is normally banded and sparsely populated. The boundary line of the shaded area in Figure 8 represents the leading non-zero entry locations in a hypothetical stiffness matrix. FEABL subprograms which operate on the master stiffness matrix incorporate logic instructions which cause the operation to be skipped if the entry K_{ij} lies in the unshaded area of Figure 8. Thus, the leading zero entries for each row are not stored, and this is where the address count vector and the LNZ vector come into play. The actual organization of the master stiffness array consists of a sequence of

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FIGURE 8

variable-length sections for the rows of the stiffness matrix.
Figure 9 illustrates a sample sequence.

FIGURE 9

The algorithms for operating on the master stiffness matrix, based on the contents of the address count and LNZ vectors are quite simple:

        (For an operation on K_{ij}: row I, column J, $J \leq I$)
        :
        :
        INDEX=ILNZ+I-1
        IF (J .LT. INTGR(INDEX)) GØ TØ 5
        KADR=IKOUNT+I-1
        KADR=INTGR(KADR)+J
        (To define address of K_{ij} in master stiffness array)
        (Operate using REAL(KADR) for K_{ij})
        :
        :
    5   (Skip operation if J<LNZ column number for the row)

2.3   Modification of the FEABL Data Storage System

The data storage system presented above has been designed for "production" computing. It has been assumed implicitly that the user will be conducting a great number of studies involving similar stress analysis problems and employing nearly the same number of degrees of freedom. Thus, the /DATA/ vector need be DIMENSIØNed only once, after which production object decks of the FEABL software may be made.

However, the FEABL software may also be placed in on-line storage in a form which will handle problems of widely varying size, with only minor modifications. These modifications are as follows:

1. Delete the CØMMØN /DATA/ REAL declaration from all programs and subprograms.

2. Delete the EQUIVALENCE (REAL(1), INTGR(1)) declaration from all subprograms. (Retain this declaration in MAIN.)

3. Use the standard declaration:
        DIMENSIØN REAL(2), INTGR(2)
    in all FEABL subroutines. (DIMENSIØN the /DATA/ vector properly in MAIN.)

4. Add the array names REAL and INTGR as arguments of

all FEABL subroutines, e.g.:

```
SUBRØUTINE ØRK(LENGTH, REAL, INTGR)
SUBRØUTINE FACTPD(REAL, INTGR)
```

When the on-line version of FEABL is used, only the DIMENSIØN
declaration for the /DATA/ vector in the user's MAIN program
need be changed to perform analyses requiring different amounts
of data storage.

SECTION 3

FEABL PROCESS CONTROL

3.1 General

For present purposes the finite element analysis of a structure will be divided into eight programming stages:

1. Establishment of input/output device codes, problem size, and address index parameters.

2. Input of the master assembly list and organization of the master stiffness matrix into corresponding segments.

3. Generation and assembly of element stiffnesses (and nodal equivalent forces, if any) in global coordinates.

4. Application of rotation transformations to the master stiffness matrix (and assembled nodal equivalent forces) at any nodes at which the boundary conditions are to be given in special coordinate systems.

5. Input of the prescribed quantities, i.e., global numbers at which displacements are to be prescribed, values of prescribed displacements, and accumulation of values of any nonzero global concentrated forces. Application of constraints to the master stiffness matrix and force/displacement vector.

6. Solution for the master displacement vector.

7. Application of inverse rotation transformations to the master displacement vector, at any nodes where a rotation was applied in stage 4, to produce a master displacement vector entirely in the global coordinate system.

8. Extraction of element displacement vectors from the global vector; calculation of element stresses from the element displacement vector.

Specific FEABL subprograms are associated with each of the above stages, according to the following table:

Stage	FEABL Subroutines	Stage	FEABL Subroutines
1	SETUP	5	BCØN
2	ØRK	6	FACTPD/FACTSD, SIMULQ
3	ASEMBL	7	RØTATE
4	RØTATE	8	XTRACT

Figure 10 illustrates the standard FEABL process sequence in terms of the eight stages described above.

Analysts who are just beginning to work with FEABL are advised to observe strictly the standard process sequence outlined above. This sequence will be described in full detail in Section 4. More experienced analysts may find it convenient to depart occasionally from the standard process in order to reduce program length.

3.2   Description of the FEABL Software

The FEABL package consists of the eight subroutines listed in the table in Subsection 3.1. Each subroutine is associated with a particular stage in the standard process sequence (except RØTATE, which is associated with both stages 4 and 7). All FEABL subprograms are ready to use for in-core finite element analysis, once the DIMENSIØN declaration for the /DATA/ vector has been inserted (see Subsections 2.2 and 2.3). The following subsections demonstrate how each of the FEABL subroutines is called, define the subroutine arguments, and outline briefly what the subroutine does. A summary table of the CØMMØN area information required by each FEABL subroutine is given in Appendix A. Listings appear in Appendix C.

User's MAIN

Stage	MAIN block	Called routines
Stage 1	Establish device codes and /SIZE/ group	SETUP
Stage 2	Input assembly list	ØRK
Stage 3	Loop over all elements	Stiffness Generator / ASEMBL
Stage 4 (Optional)	Rotations: One CALL per pair or triplet of DOF at each node required	RØTATE
Stage 5	Input boundary condition information	BCØN
Stage 6		FACTPD or FACTSD / SIMULQ
Stage 7 (if Stage 4 appears)	Inverse rotations	RØTATE
Stage 8	Loop over all elements	XTRACT / Stress Generator

FIGURE 10

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3.2.1 Housekeeping Setup Subroutine for /DATA/ Vector (SETUP)

CALL SETUP(LENGTH, NCØN, MASTRL)

LENGTH – A scalar integer numerically equal to the dimension which the user has assigned to the /DATA/ vector.

NCON – A scalar integer greater than or equal to the total number of degrees of freedom, in the assembled structure, at which displacements are to be prescribed.

MASTRL – A scalar integer greater than or equal to the total number of words required for the master assembly list.

Based on NCØN, MASTRL and the total number of degrees of freedom in the entire assembled structure (NDT, in the /SIZE/ group), subroutine SETUP organizes the /DATA/ vector by calculating the address index parameters in the /BEGIN/ and /END/ groups, except for the index LK which defines the end of the block reserved for the master stiffness matrix. SETUP uses the argument LENGTH to test whether the user's /DATA/ vector has at least enough storage available to accommodate the first four data blocks (constraint vector, address count vector, LNZ vector and master assembly list). If the /DATA/ vector is too short, SETUP estimates the total length required for all six data blocks, based on a reasonable population density for the lower triangle of the master stiffness matrix, prints the estimate and aborts the run. If the first four blocks can be accommodated, the constraint vector is filled with zeros and control is returned to MAIN.

3.2.2 Subroutine for Detailed Organization of the Master Stiffness Matrix Block (ØRK)

CALL ORK(LENGTH)

LENGTH - A scalar integer numerically equal to the dimension which the user has assigned to the /DATA/ vector.

Subroutine ØRK produces, in essence, a map for the master stiffness matrix like that of Figure 9, using the information contained in the master assembly list to calculate the correct values of the entries in the LNZ vector. This is accomplished by first setting each LNZ column number equal to its row number (diagonal matrix), and then examining the assembly information element by element to re-set the LNZ column numbers, according to the following algorithm:

1. The smallest global number N for the element is found.
2. The element's global numbers are then treated as row numbers. If the LNZ column number corresponding to a row (global number) is greater than N, its value is re-set to N.

Once the LNZ vector has been established, subroutine ØRK uses an accumulation process to calculate the absolute address of the diagonal entry of each row. By convention, the diagonal is the only entry stored for the first row and is therefore located by the address index parameter IK:

INTGR(IKØUNT)=IK

Subsequent entries are located by the algorithm:

INTGR(IKØUNT+M-1)=INTGR(IKØUNT+M-2)+M+1-INTGR(ILNZ+M-1)
(Diag Addr for Mth row)=(Diag Addr for row M-1)+(Total no. of nonzero entries in Mth row)

Conveniently, the address of the diagonal for the last row in the master stiffness matrix is also the correct value of LK. Subroutine ØRK now tests the /DATA/ vector by means of the argument LENGTH. If the /DATA/ vector is too short for the problem data, an exact calculation of its required length is

output and the run is aborted. If sufficient storage is available, the stiffness matrix map is output and the address count vector is modified by:

$$INTGR(IK\text{\O}UNT+M-1)=INTGR(IK\text{\O}UNT+M-1)-M; \quad 1 \leq M \leq NDT$$

for all rows, M. This saves repeated subtraction of the row number in later subprograms. The correct algorithm for locating K_{MN} in the /DATA/ vector is now:

$$KADR=INTGR(IK\text{\O}UNT+M-1)+N; \quad M \geq N$$
K_{MN} is assigned to REAL(KADR)

Since the next stage of the analysis will involve the accumulation of data in the master stiffness matrix (and perhaps in the force/displacement vector), $\text{\O}RK$'s last action before returning control to MAIN is to fill these two data blocks with floating point zeros.

3.2.3 Element Assembly Subroutine (ASEMBL)

CALL ASEMBL(LNUM, NDE, ELK, ELQ)

LNUM – A positive scalar integer = user's global element number

NDE   – A scalar integer = total number of degrees of freedom possessed by the element which is about to be assembled

ELK   – A floating point, two-dimensional square array which contains the stiffness matrix of the element about to be assembled.

ELQ   – A floating point vector which contains the equivalent nodal forces for the element about to be assembled, or which contains floating point zeros if there are no equivalent nodal forces.

Since ELK and ELQ are variably DIMENSI$\text{\O}$Ned in this subprogram, elements having different numbers of degrees of freedom can be handled automatically. However, these arguments must be DIMENSI$\text{\O}$Ned explicitly in the user's MAIN program. For example, suppose a structure is to be analyzed in plane stress with a

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combination of 3-node triangle elements (6 degrees of freedom) and 4-node rectangle elements (8 degrees of freedom). Then the declaration:

DIMENSIØN TRIK(6,6), TRIQ(6), RECK(8,8), RECQ(8)

might appear in MAIN. A triangle element would be assembled by the instruction:

CALL ASEMBL(LNUM, 6, TRIK, TRIQ)

while a rectangle element would be assembled by:

CALL ASEMBL (LNUM, 8, RECK, RECQ)

Subroutine ASEMBL can handle elements having as many as one hundred degrees of freedom (not a serious restriction). If the assembly of larger elements is attempted, ASEMBL will abort the run and tell the user to change the DIMENSIØN of one of its internal parameters.

ASEMBL examines the section of the master assembly list belonging to element number LNUM and records the values of the global displacement numbers in an internal vector called MNUM. Then, each entry K_{ij} of the lower triangle of the element stiffness matrix ($i \geq j$), is accumulated into its proper place in the master stiffness matrix, according to the algorithm:

I=MNUM(i)

J=MNUM(j)

$K_{ij} \rightarrow K_{IJ}$ for $I \geq J$

$K_{ij} \rightarrow K_{JI}$ for $I < J$

(See Subsection 3.2.1 for address algorithm for K_{IJ}.) ASEMBL does not use the upper triangle ($i < j$) of the element stiffness matrix; the user may omit calculating these entries to save time. The vector ELQ of element equivalent nodal forces is accumulated in the same manner into the proper locations in the master force vector (fifth block of the /DATA/ vector).

3.2.4 Rotation Transformation Subroutine (RØTATE)

CALL RØTATE(NØDE, IRØW, JRØW, KRØW, ZANGLE, YANGLE, XANGLE)

NØDE – A scalar integer "convenience" number for the user. May be positive or negative but not zero. (See explanation below.)

IRØW ⎫ Three scalar integers equal to the global numbers
JRØW ⎬ of the degrees of freedom at the node at which
KRØW ⎭ the rotation is to be performed.

ZANGLE ⎫ Floating point values of the three Euler rotation
YANGLE ⎬ angles in degree measure.
XANGLE ⎭

The argument NØDE does not enter directly into the transformation calculations, but is printed out in the heading of the information supplied by RØTATE. Normally, the user assigns either his node numbers or the series 1, 2, 3,... to a set of rotations. A positive value of NØDE causes the full subprogram to execute, and is used for stage 4 in the process sequence. In stage 7, only the solution displacement vector requires transformation; a negative value of NØDE will cause execution of the now unwanted transformation of the stiffness matrix to be skipped.

The global numbers of the degrees of freedom at the node being rotated may be contiguous (e.g., 13, 14, 15) or separated (e.g., 5, 10, 15) depending upon the global numbering scheme the user has adopted. Either form is acceptable to this subprogram. If only two degrees of freedom are to participate in the rotation (as, for example, in plane stress problems), an integer zero must be specified for the third global number. If six degrees of freedom are to participate (e.g., shell elements), two separate rotations are required, i.e., one CALL for the translational and one for the rotational degrees of freedom. Errors in the global number arguments (e.g., repeating a number or too many zeros) will cause an abort.

The three Euler angle arguments must be specified according to the conventions illustrated in Figure 11. Let XYZ be the user's global cartesian axis system, with respect to which the element stiffness and equivalent nodal forces have been generated and assembled. Let $\tilde{X}\tilde{Y}\tilde{Z}$ by a local coordinate system with respect to which the displacement constraints are to be specified.

FIGURE 11

The following conventions have been adopted in subroutine RØTATE:

1. Axes X, Y are first rotated through angle ZANGLE about axis Z to the intermediate positions x, y.

2. Axes x, Z are then rotated through angle YANGLE about axis y, x to the final position $\tilde{X}$ and Z to an intermediate position z.

3. Axes y, z are then rotated through angle XANGLE about axis $\tilde{X}$ to their final positions $\tilde{Y}$, $\tilde{Z}$.

4. Positive angles obey the "right hand rule" of vector analysis.

Subroutine ROTATE forms the matrix of direction cosines:

$$
\mathbf{D} = \begin{bmatrix} \cos(\tilde{X}, X) & \cos(\tilde{X}, Y) & \cos(\tilde{X}, Z) \\ \cos(\tilde{Y}, X) & \cos(\tilde{Y}, Y) & \cos(\tilde{Y}, Z) \\ \cos(\tilde{Z}, X) & \cos(\tilde{Z}, Y) & \cos(\tilde{Z}, Z) \end{bmatrix}
$$

from the values of the Euler angles supplied in the argument list. Then, if argument NØDE is positive, RØTATE applies the following transformations to portions of the master stiffness matrix **K** where i, j, k are the global numbers specified by the user:

$$
\mathbf{D} \begin{bmatrix} K_{ii} & K_{ij} & K_{ik} \\ K_{ji} & K_{jj} & K_{jk} \\ K_{ki} & K_{kj} & K_{kk} \end{bmatrix} \mathbf{D}^T
$$

$$
\mathbf{D} \begin{bmatrix} K_{i1} & K_{12}\cdots K_{1,i-1} & K_{1,i+1}\cdots K_{1,j-1} & K_{1,j+1}\cdots K_{1,k-1} & K_{1,k+1}\cdots \\ K_{j1} & K_{j2}\cdots K_{j,i-1}\cdots \\ K_{k1} & K_{k2}\cdots \end{bmatrix}
$$

$$
\begin{bmatrix}
K_{11} & K_{1j} & K_{1k} \\
K_{21} & K_{2j} & K_{2k} \\
\vdots & \vdots & \vdots \\
K_{i-1,i} & K_{i-1,j} & \\
K_{i+1,i} & & \\
\vdots & & \\
K_{j-1,i} & & \\
K_{j+1,i} & & \\
\vdots & & \\
K_{k-1,i} & & \\
K_{k+1,i} & & \\
\vdots & &
\end{bmatrix} \mathbf{D}^{\mathsf{T}}
$$

Finally, the corresponding entries of the force/displacement vector are transformed according to:

$$
\mathbf{D}
\begin{Bmatrix}
Q_i \\
Q_j \\
Q_k
\end{Bmatrix}
$$

A few examples will illustrate the proper use of subroutine RØTATE. First, suppose that global degrees of freedom 1 and 2 are to be rotated by 45 degrees at node 1 in a plane elasticity problem. The correct instruction is:

    CALL RØTATE(1, 1, 2, 0, 45., 0., 0.)

To return the force/displacement vector to the user's global axis system after the displacement solution has been obtained:

    CALL RØTATE(-1, 1, 2, 0, -45., 0., 0.)

Note that only the first angular argument ZANGLE is used in two-dimensional problems. In a three-dimensional problem, to rotate

degrees 4, 5, 6 at node 2 through angles of 15, 30 and 45 degrees about the Z, y and $\tilde{X}$ axes:

      CALL RØTATE(2, 4, 5, 6, 15., 30., 45.)

and after the displacement solution has been obtained:

      CALL RØTATE(-2, 4, 5, 6, -45., -30., -15.)

Note that not only the signs, but also the <u>order of application</u> of the angles is reversed. However, some care is required in three dimensions. Suppose degrees 7, 8, 9 at node 3 were rotated only about axes Z and y:

      CALL RØTATE(3, 7, 8, 9, 20., 40., 0.)

Then to reverse the rotation after obtaining the displacement solution:

      CALL RØTATE(-3, 7, 8, 9, -40., -20., 0.)

Note that the unused XANGLE does <u>not</u> participate in the reversal. The instruction:

      CALL RØTATE(-3, 7, 8, 9, 0., -40., -20.)

would be incorrect.

3.2.5   Boundary Constraint Subroutine (BCØN)

CALL BCØN

Let u, Q represent respectively the global[*] displacement vector and the global force vector. The result of the first four program stages has been to supply a right-hand side and the stiffness coefficients K_{ij} for the force-displacement relations:

$$Ku = Q$$

However, the value of u is known at some degrees of freedom, with the corresponding Q unknown, while the value of Q is known

---

[*]The term "global" should be taken in a more general sense here. It refers to the final set of coordinate systems $\tilde{X}\tilde{Y}\tilde{Z}$ in which the user will prescribe his boundary conditions.

at the other degrees.

Conceptually, the set of all degrees of freedom in the structure may be divided into two subsets: Those at which forces are prescribed (F) and those at which displacements are prescribed (D). When the user specifies the values of the prescribed quantities, Q becomes a "force/displacement" vector in fact:

$$Q = \left\{ \begin{array}{c} \hat{Q}_F \\ \hat{u}_D \end{array} \right\} \quad \begin{array}{l} (\text{"}\wedge\text{" means a} \\ \text{prescribed quantity}) \end{array}$$

and:

u_F = unknown displacements
Q_D = unknown reaction forces

The force-displacement relations may be partitioned in a similar manner:

$$\begin{bmatrix} K_{FF} & K_{FD} \\ K_{FD}^T & K_{DD} \end{bmatrix} \left\{ \begin{array}{c} u_F \\ \hat{u}_D \end{array} \right\} = \left\{ \begin{array}{c} \hat{Q}_F \\ Q_D \end{array} \right\}$$

Subroutine BCØN transforms the force-displacement relations from the above form to:

$$\begin{bmatrix} K_{FF} & 0 \\ 0 & I \end{bmatrix} \left\{ \begin{array}{c} u_F \\ \hat{u}_D \end{array} \right\} = \left\{ \begin{array}{c} \hat{Q}_F - K_{FD}\hat{u}_D \\ \hat{u}_D \end{array} \right\}$$

With the above "right hand side" in the force/displacement vector, standard equation-solving techniques may be used to produce the solution displacement vector. Before making the transformation, BCØN arranges the entries of the constraint vector (Subsection 2.2.1)

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in ascending order and checks for the presence of global numbers
(positive integers). The run is aborted if no global numbers
are found.

3.2.6 Stiffness Matrix Factoring Subroutine (FACTPD/FACTSD)

CALL FACTPD

or

CALL FACTSD

The force-displacement relations are solved by Choleski's
direct method, which consists of two programming steps:

1. The master stiffness matrix is factored into a triple
   product.

2. The displacements are solved for sequentially, in
   three sub-steps.

Subroutine FACTPD/FACTSD accomplishes the first programming step.
The master stiffness matrix K is factored into the form:

$$K = LDL^T$$

where L is a lower triangular matrix:

$$L = \begin{bmatrix} 1 & 0 & 0 & 0\ldots \\ L_{21} & 1 & 0 & 0\ldots \\ L_{31} & L_{32} & 1 & 0\ldots \\ L_{41} & L_{42} & L_{43} & 1\ldots \\ \vdots & \vdots & \vdots & \vdots \end{bmatrix}$$

and $D = \begin{bmatrix} D_1 & D_2 \ldots \end{bmatrix}$ is a diagonal matrix. The factor-
ing algorithms are:

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$$L_{mn} = \frac{1}{D_n} \left[K_{mn} - \sum_{j=J(m,n)}^{n-1} \left(D_j L_{mj} L_{nj} \right) \right] ; \; n = J(n), \ldots, m-1 \Bigg\} \; m = 1, 2, \ldots, NDT$$

$$D_m = K_{mm} - \sum_{j=J(m)}^{m-1} \left(D_j L_{mj}^2 \right)$$

where:

J(n) = Leading Nonzero Entry Column Number for Row n
J(m) = Leading Nonzero Entry Column Number for Row m
J(m,n) = max(J(n), J(m))

The master stiffness matrix is destroyed and replaced by the entries of **L** and **D** as the factoring process is executed. The entries L_{mn} (m>n) and D_m are stored respectively at K_{mn} (m>n) and K_{mm}. The unit diagonal entries of **L** are not stored.

FACTPD/FACTSD tests the entries of **D** for nonsingularity and positive-definiteness as they are created. **K** is positive-definite if all $D_m > 0$. Rows for which $D_m < 0$ are reported. If any D_m is found to equal zero exactly, **K** is singular; the row in which the singularity was discovered is reported and the run is aborted.

The names FACTPD and FACTSD refer to different entry points in this subprogram. If **K** is supposed to be positive-definite (as in the case of a structure analyzed by compatible displacement elements), FACTPD should be called. If errors were made in the assembly or constraint stages of the program, they will appear now as $D_m < 0$ in one or more rows, and FACTPD will abort the run. If **K** is not necessarily positive-definite (as in the case of a structure analyzed by hybrid stress-displacement elements), FACTSD should be called. FACTSD continues execution even if there exist $D_m < 0$. If the FEABL software is to be converted to the on-line storage version (Section 2.3), be sure to make the array names REAL, INTGR arguments of both entry points.

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FACTPD/FACTSD also makes a rough estimate of the conditioning of $\mathbf{K}$ by calculating the so-called rounding error parameter (see Ref. 2, pg. 81):

$$E = \min\left(|D_m| / |K_{mm}| \right)$$

E is a measure of how many significant figures of information have been lost in the diagonal entries, as a result of the factoring algorithm. The run is aborted if $E < 10^{-5}$. The user must keep in mind the fact that the rounding error parameter is an imperfect conditioning measure. If FACTPD/FACTSD reports that no significant figures have been lost, it does not necessarily follow that there is no error in the displacement solution. Advanced techniques for realistic solution error estimates are discussed in Section 5.

3.2.7 Subroutine for Solution of Simultaneous Equations (SIMULQ)

CALL SIMULQ(ENERGY)

ENERGY - A scalar floating point variable, the value of which is undefined when SIMULQ is called.

With the stiffness matrix in factored form $\mathbf{K} = \mathbf{L}\,\mathbf{D}\mathbf{L}^T$, let $\mathbf{P} = \mathbf{L}^T\boldsymbol{u}$ and $\mathbf{R} = \mathbf{D}\mathbf{P}$. Then:

$$\mathbf{L}\mathbf{R} = \mathbf{Q}$$

can be solved sequentially for $R_1, R_2, \ldots, R_{NDT}$ via the algorithm:

$$R_1 = Q_1$$

$$R_m = Q_m - \sum_{j=J(m)}^{m-1} \left(L_{mj} R_j \right) \; ; \; m = 2, 3, \ldots, NDT$$

Then:

$$P_m = R_m / D_m \; ; \; m = 1, 2, \ldots, NDT$$

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Finally, the displacement solution is obtained by solving $\mathbf{L}^T \mathbf{u} = \mathbf{P}$ sequentially for $u_{NDT}, u_{NDT-1}, \ldots, u_1$ according to:

$$u_{NDT} = P_{NDT}$$

$$u_m = P_m - \sum_{j=m+1}^{NDT} (L_{jm} u_j) \; ; \; m = NDT-1, NDT-2, \ldots, 1$$

As SIMULQ carries out these three sub-steps, the prescribed vector $\mathbf{Q}$ is first replaced by $\mathbf{R}$, then $\mathbf{R}$ is replaced by $\mathbf{P}$ and finally $\mathbf{P}$ is replaced by $\mathbf{u}$. $\mathbf{Q}$ and $\mathbf{u}$ are printed out by SIMULQ.

Since the approximate value of the strain energy in the structure is often useful to the analyst, this quantity is calculated by SIMULQ during execution of the solution steps. If the factored form of $\mathbf{K}$ is introduced into the strain energy expression, there results:

$$\text{Strain Energy} = \tfrac{1}{2} \mathbf{u}^T \mathbf{K} \mathbf{u} = \tfrac{1}{2} (\mathbf{u}^T \mathbf{L}) \mathbf{D} (\mathbf{L} \mathbf{u}) = \tfrac{1}{2} \mathbf{P}^T \mathbf{D} \mathbf{P}$$

Thus, the straightforward algorithm:

$$\text{Strain Energy} = \tfrac{1}{2} \sum_{j=1}^{NDT} D_j (P_j)^2$$

can be used. The value of the strain energy thus calculated is printed out. At the end of execution of subroutine SIMULQ, the force/displacement vector (fifth block in the /DATA/ vector) contains the master displacement solution vector, and the strain energy value has been assigned to the argument ENERGY.

3.2.8 Subroutine for Extraction of Element Displacements from the Global Displacement Vector (XTRACT)

CALL XTRACT(LNUM, NDE, ELQ)

LNUM – A positive scalar integer equal to user's global element number

NDE - A scalar integer equal to the total number of degrees
of freedom in the element

ELQ - A floating point vector at least NDE words long.
The entries of ELQ are undefined when XTRACT is
called.

If the analyst desires to calculate stress or strain dis-
tributions in his structure, he commonly uses transformations
between stress or strain and nodal displacements:

$$\varepsilon = B\,u_{el} \qquad\qquad \sigma = E\varepsilon$$

where ε, σ are respectively vectors of strain and stress com-
ponents at selected points in the element domain, and where
u_{el} is the vector of element nodal displacements, i.e., a subset
of the global vector u . Subroutine XTRACT selects the correct
subset u_{el} out of u , based upon the information contained
in the user's master assembly list. The values of u_{el} are
placed in the argument vector ELQ.

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A SAMPLE FEABL PROGRAM

In order to illustrate further the data location and process sequence interfacing, sample user programs will be developed for analysis of the truss structure shown in Figure 12. The structure consists of 16 bars and 18 degrees of freedom in the XY plane.

FIGURE 12

Local number conventions for the ⓝ th typical bar element are shown at the right. With these conventions, the stiffness matrix of the typical bar element is given by:

$$\mathbf{k}_{\circledR} = \frac{EA}{\ell} \begin{bmatrix} \cos^2\alpha & & & (Symmetric) \\ \sin\alpha\cos\alpha & \sin^2\alpha & & \\ -\cos^2\alpha & -\sin\alpha\cos\alpha & \cos^2\alpha & \\ -\sin\alpha\cos\alpha & -\sin^2\alpha & \sin\alpha\cos\alpha & \sin^2\alpha \end{bmatrix}$$

where:

$$\ell = \sqrt{(X_2-X_1)^2 + (Y_2-Y_1)^2}$$

$$\cos\alpha = \frac{1}{\ell}(X_2-X_1)$$

$$\sin\alpha = \frac{1}{\ell}(Y_2-Y_1)$$

and where E, A are the bar's modulus and cross section area. Also, once the displacements u_1, u_2, u_3, u_4 are known for the element, its elongation may be calculated as:

$$\delta = (u_3-u_1)\cos\alpha + (u_4-u_2)\sin\alpha$$

and the load in the bar is then given by:

$$P = EA\delta/\ell$$

where ℓ and α are defined as above.

The user decides to read the properties and nodal coordinates for each element, each time they are required (a rather inefficient procedure). Reading is to be done by the generator subroutine, rather than in MAIN; however, the card reader device code will be established in MAIN. Therefore, the user programs his stiffness matrix and stress generator subroutines as follows:

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Stiffness Matrix Generator

```
        SUBRØUTINE BARK (ELK, ELQ)
        DIMENSIØN ELK (4, 4), ELQ (4)
        CØMMØN /IØ/ KR, KW, KP, KT1, KT2, KT3
C  INTERFACE WITH IØ CØNTRØL PARAMETERS IS ØPTIØNAL
  91    FØRMAT (6E10.3)
        READ (KR, 91) X1, Y1, X2, Y2, E, A
C  CALCULATE BAR LENGTH
        BARL=SQRT ((X2-X1)** 2+(Y2-Y1)** 2)
C  CALCULATE SINE AND CØSINE
        S=(Y2-Y1)/BARL
        C=(X2-X1)/BARL
C  CALCULATE ENTRIES IN LØWER TRIANGLE ØF ELK-ASEMBL DØES
C  NØT USE UPPER TRIANGLE
        ELK (1, 1)=E*A*C*C/BARL
        ELK (2, 1)=E*A*S*C/BARL
        .
        .
        .
        etc.
        .
        .
        ELK (4, 4)=E*A*S*S/BARL
C  ESTABLISH ZERØ ELEMENT EQUIVALENT NØDAL FØRCES
        DØ 10 I=1, 4
  10    ELQ (I)=0.
        RETURN
        END
```

Stress Generator

```
        SUBRØUTINE BARF (LNUM, ELQ)
        DIMENSIØN ELQ (4)
        CØMMØN /IØ/ KR, KW, KP, KT1, KT2, KT3
  91    FØRMAT (6E10.3)
  92    FØRMAT (21HOBAR FØRCE IN BAR NØ., I4, 2H =, E10.3, 3H LB)
        READ (KR, 91) X1, Y1, X2, Y2, E, A
        BARL = SQRT ((X2-X1)** 2+(Y2-Y1)**2)
        S=(Y2-Y1)/BARL
        C=(X2-X1)/BARL
        FORCE=E*A*(C*(ELQ(3) - ELQ (1)) + S* (ELQ(4)-ELQ (2)))/BARL
        WRITE (KW, 92) LNUM, FØRCE
        RETURN
        END
```

Subroutines BARK and BARF are ready to use in conjunction with
the FEABL software. The user now begins the construction of
his MAIN program, one stage at a time.

<u>Stage 1: Program Heading (Data Location Interface), Device</u>
<u>Code and Problem Size Establishment</u>

The user estimates that a 1000-word /DATA/ vector will
be more than adequate for the problem. Required device codes
are the card reader (5) and printer (6).

```
C   MAIN PRØGRAM FØR SØLUTIØN ØF TRUSS PRØBLEM
      DIMENSIØN REAL(1000), INTGR(1000)
      DIMENSIØN ELK(4, 4), ELQ(4)
      CØMMØN /IØ/ KR, KW, KP, KT1, KT2, KT3
      CØMMØN /SIZE/ NET, NDT
      CØMMØN /BEGIN/ ICØN, IKØUNT, ILNZ, IMASTR, IQ, IK
      CØMMØN /END/ LCØN, LKØUNT, LLNZ, LMASTR, LQ, LK
      EQUIVALENCE (REAL(1), INTGR(1))
      KR=5
      KW=6
      NET=16
      NDT=18
      CALL SETUP (1000, 5, 80)
C   END ØF STAGE 1
```

The user has called for space for five constraints: master
displacements 2 (after a -45° rotation of 1 and 2), 7, 8, 13
and 14. The assembly list must allow room for 16 elements x
(4 DOF plus 1 pointer per element) = 80 words. The first
declaration:

```
      DIMENSIØN REAL(1000), INTGR(1000)
```

is also duplicated and placed in each FEABL subroutine.

<u>Stage 2: Assembly List Input and Organization of **K**</u>

The user recognizes that the pointers will occupy locations
IMASTR to IMASTR+15 (see Section 2.2.4). He chooses to write
specific assignment instructions for each element:

```
      INTGR(IMASTR)=IMASTR+16 (Pointer for 1st element)
      INTGR(IMASTR+16)=1
      INTGR(IMASTR+17)=2
      INTGR(IMASTR+18)=7
      INTGR(IMASTR+19)=8
      .
      .
      .
```

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Stage 1: <u>Program Heading (Data Location Interface), Device</u>
<u>Code and Problem Size Establishment</u>

The user estimates that a 1000-word /DATA/ vector will
be more than adequate for the problem. Required device codes
are the card reader (5) and printer (6).

```
C  MAIN PRØGRAM FØR SØLUTIØN ØF TRUSS PRØBLEM
       DIMENSIØN REAL(1000), INTGR(1000)
       DIMENSIØN ELK(4, 4), ELQ(4)
       CØMMØN /IØ/ KR, KW, KP, KT1, KT2, KT3
       CØMMØN /SIZE/ NET, NDT
       CØMMØN /BEGIN/ ICØN, IKØUNT, ILNZ, IMASTR, IQ, IK
       CØMMØN /END/ LCØN, LKØUNT, LLNZ, LMASTR, LQ, LK
       EQUIVALENCE (REAL(1), INTGR(1))
       KR=5
       KW=6
       NET=16
       NDT=18
       CALL SETUP (1000, 5, 80)
C  END ØF STAGE 1
```

The user has called for space for five constraints: master
displacements 2 (after a -45° rotation of 1 and 2), 7, 8, 13
and 14. The assembly list must allow room for 16 elements x
(4 DOF plus 1 pointer per element) = 80 words. The first
declaration:

```
       DIMENSIØN REAL(1000), INTGR(1000)
```

is also duplicated and placed in each FEABL subroutine.

Stage 2: <u>Assembly List Input and Organization of **K**</u>

The user recognizes that the pointers will occupy locations
IMASTR to IMASTR+15 (see Section 2.2.4). He chooses to write
specific assignment instructions for each element:

```
       INTGR(IMASTR)=IMASTR+16 (Pointer for 1st element)
       INTGR(IMASTR+16)=1
       INTGR(IMASTR+17)=2
       INTGR(IMASTR+18)=7
       INTGR(IMASTR+19)=8
            ⋮
            ⋮
```

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```
          .
          .
          .
      INTGR(IMASTR+7)=IMASTR+44 (pointer for 8th element)
      INTGR(IMASTR+44)=7
      INTGR(IMASTR+45)=8
      INTGR(IMASTR+46)=9
      INTGR(IMASTR+47)=10
          .
          .
          .
      INTGR(IMASTR+15)=IMASTR+76 (pointer for 16th element)
      INTGR(IMASTR+76)=7
      INTGR(IMASTR+77)=8
      INTGR(IMASTR+78)=15
      INTGR(IMASTR+79)=16 (80th location in master assembly list)
```

At this point, it is advisable to dump the assembly list for
debugging purposes if there is any doubt about its accuracy.
Dumping may be done by:

```
      WRITE (KW, 5) (INTGR(I), I=IMASTR, LMASTR)
   5  FØRMAT (1X, 10I10)
```

Finally:

```
      CALL ØRK(1000)
   C  END ØF STAGE 2
```

Stage 3: Generation and Assembly of Element Properties

In this stage, the user merely invokes the appropriate
subroutines in a loop.

```
      DØ 10 LNUM=1, NET
      CALL BARK (ELK, ELQ)
      CALL ASEMBL (LNUM, 4, ELK, ELQ)
  10  CØNTINUE
   C  END ØF STAGE 3
```

Stage 4: Rotation Transformations

In the present problem, only the lower left node (Figure 12)
requires rotation. Master degrees of freedom number 1 and 2
must be rotated by -45 degrees.

```
      CALL ROTATE(1, 1, 2, 0, -45., 0., 0.)
                           ↑      ↑   ↑
                           Quantities for 3-D problems are not used
   C  END ØF STAGE 4
```

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Stage 5: Boundary Conditions

Displacements 2, 7, 8, 13, 14 are constrained, and the constraint vector starts in INTGR(1) (ICØN=1). Therefore:

```
INTGR(1)=2
INTGR(2)=7
INTGR(3)=8
INTGR(4)=13
INTGR(5)=14
```

The constrained displacements are all prescribed to be zero. Only the nonzero prescribed forces need be input; these are 1,000 lb. each at the 15th and 17th degrees of freedom. Therefore:

```
C   PRESCRIBED DISPLACEMENTS
    REAL(IQ+1)=0.
    REAL(IQ+6)=0.
    REAL(IQ+7)=0.
    REAL(IQ+12)=0.
    REAL(IQ+13)=0.
C   ACCUMULATE PRESCRIBED FØRCES
    REAL(IQ+14)=REAL(IQ+14)+1000.
    REAL(IQ+16)=REAL(IQ+16)+1000.
    CALL BCØN
C   END ØF STAGE 5
```

Stage 6: Choleski Solution

```
    CALL FACTPD
    CALL SIMULQ(ENERGY)
C   END ØF STAGE 6
```

Stage 7: Inverse Rotation to Obtain Global Displacements in Global Cartesian Coordinates

A rotation of -45° was performed at node 1 before the boundary conditions were imposed. Since all element calculations are done with respect to the unrotated XY axis system, this rotation must be reversed before element stresses are calculated:

```
    CALL RØTATE(-1, 1, 2, 0, 45., 0., 0.)
C   END ØF STAGE 7
```

Stage 8: Calculation of Element Stresses

In this case, bar forces are to be calculated. Again, the user merely invokes the proper subroutines:

```
        DØ 20 LNUM=1, NET
        CALL XTRACT (LNUM, 4, ELQ)
        CALL BARF (LNUM, ELQ)
     20 CØNTINUE
C       END ØF PRØG
        STØP
        END
```

If the reader has grasped the material presented up to this point, he now has enough familiarity with FEABL to use it, albeit somewhat inefficiently. However, some additional degree of sophistication is desirable for the instructions which input problem data such as the assembly list. It is apparent that a straightforward set of instructions such as that given in the sample program is quite cumbersome, especially for problems involving more elements with more degrees of freedom per element. Improved programming techniques are discussed in the next section.

SECTION 5

ADVANCED TECHNIQUES WITH FEABL

5.1   Efficient Programming for Large Problems

When a structure is to be analyzed with a large number
of elements and many degrees of freedom, there often occurs a
definite trend toward <u>regularity</u> in the element set. The analyst
should recognize two distinct forms of regularity, and he should
be prepared to take advantage of each in writing his MAIN program.

First, consider the truss structure shown in Figure 13,
with all bays having the same dimension vertically and horizontally.

FIGURE 13

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The bar elements may be said to possess a degree of regularity, in that there are only four different types of elements in the structure:

1. Horizontal bars
2. Vertical bars.
3. +45° diagonals.
4. -45° diagonals.

Since an element stiffness matrix depends upon nodal coordinates only through differences in the coordinate values (see beginning of Section 4), there will occur only four independent element stiffness matrices for the structure of Figure 13: one for each of the element types listed above. Thus, for example, the element stiffness matrix for a typical horizontal bar may be generated once and assembled repeatedly (32 times for the present case). The technique may be repeated again for the 30 vertical bars, 24 +45° diagonals and 24 -45° diagonals, achieving considerable savings in execution time.

The second form of regularity involves the way in which displacements and elements may be numbered. Often an element set is topologically equivalent to a rectangular or square array, even if it is not geometrically regular. For example, suppose the trapezoidal continuum shown in Figure 14(a) is to be analyzed in plane stress using 32 3-node triangle elements. For numbering purposes the element and displacement sets are topologically equivalent to the square net shown in Figure 14(b). It is then possible to generate the assembly list for the structure by a double DØ loop. Each element is "located" in the structure via an intersection of one of the element strings LX and one of the element strings LY. The element number and its master displacement numbers may be generated from the values of LX and LY. Taking advantage of numbering regularity enables the user to input large sections of the assembly list with relatively few FØRTRAN instructions.

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FIGURE 14

5.2   A Numbering Strategy for Planar Problems

Suppose the set of MxN rectangular plane stress el?ments shown in Figure 15 is to be assembled, where M, N may be quite large. If the numbering strategy shown in the figure is adopted, beginning with element 1 and displacements 1, 2 at the lower left corner and ending with element MN and displacements 2(M+1) (N+1)-1, 2(M+1) (N+1) at the upper right corner, then each element number can be given as a function of its LY, LX string coordinates:

LNUM=LY+N*(LX-1) where 1≤LY≤N and 1≤LX≤M

In order to create similar functions for the element assembly list, a local numbering convention must be adopted. If the

FIGURE 15

convention shown is Figure 15 is used, the global DOF numbers
may be calculated as follows: Local displacements 1, 2 can be
considered to lie on strings LX, LY; hence:

    (Local #1)=2*LY-1+2*(N+1)*(LX-1)
    (Local #2)=(Local #1)+1

Local displacements 7, 8 follow directly from the master numbering
scheme:

    (Local #7)=(Local #2)+1=(Local #1)+2
    (Local #8)=(Local #7)+1=(Local #1)+3

Local displacements 3, 4 are shifted one LX string to the right;
hence they are 2*(N+1) ahead of 1, 2:

    (Local #3)=(Local #1)+2*(N+1)
    (Local #4)=(Local #1)+2*(N+1)+1=(Local #2)+2*(N+1)

and finally:

    (Local #5)=(Local #1)+2*(N+1)+2=(Local #7)+2*(N+1)
    (Local #6)=(Local #1)+2*(N+1)+3=(Local #8)+2*(N+1)

    With the algorithms derived above, it is quite easy to
develop an efficient procedure to input the entire assembly list
into the /DATA/ vector. The reader will recall that the pointer

for element number LNUM is stored in INTGR(IMASTR+LNUM-1).
One additional integer variable, called NEXT, is required. NEXT
is to be incremented after completion of the input for one ele-
ment, so that the next value of NEXT is the address of the next
available location in the master assembly list, i.e., NEXT=
the value of the pointer for the next element. The entire assembly
list is then generated and input by the following set of FØRTRAN
instructions:

```
        (Values of M and N defined elsewhere in program)
        :
        :
  C   INITIALIZE PØINTER VALUE
      NEXT=IMASTR+NET
  C   LØØP ØVER ELEMENT STRINGS
      DØ 20 LX=1, M
      DØ 20 LY=1, N
      LNUM=LY+N*(LX-1)
      IPTR=IMASTR+LNUM-1
  C   ESTABLISH PØINTER FØR THE ELEMENT
      INTGR(IPTR)=NEXT
  C   CALCULATE JD1=1 LESS THAN MASTER NØ. ØF 1ST DØF, JD3=1
  C   LESS THAN MASTER NØ. ØF 3RD DØF
      JD1=2*(LY-1)+2*(N+1)*(LX-1)
      JD3=JD1+2*(N+1)
  C   ASSIGN DOF MASTER NOS. TO LOCATIONS NEXT, NEXT + 1,..., NEXT + 7
  C   IN /DATA/ VECTØR
      J=0
      K=0
      DØ 10 I=1,8
      INDEX=NEXT+I-1
      IF(I .GT. 2 .AND. I .LT. 7) GØ TØ 5
      J=J+1
      JD=JD1
      II=J
      GØ TØ 10
   5  JD=JD3
      K=K+1
      II=K
  10  INTGR(INDEX)=JD+II
  C   INCREMENT NEXT LØCATIØN
  20  NEXT=NEXT+8
  C   STØRE A ZERØ IN NEXT LØCATIØN (ASSUMING EXCESS STØRAGE
  C   IN MASTER ARRAY)
      IF (NEXT .LE. LMASTR) INTGR(NEXT)=0
        :
        :
```

The last IF statement above is a good form of insurance for cases in which the user may have overestimated the number of words required for his master assembly list.

A similar scheme for structures consisting of triangle elements (Figure 14) can be developed if the "right-side-up" and "upside-down" elements are treated as separate sets.

The automatic generation technique derived above was developed for a structure with complete topological regularity; however, it may be extended to cover large portions of less regular structures with relatively little additional programming. There is also a hidden advantage: the numbering strategy adopted in Figure 15 is not only easy to produce, but also minimizes the population of the master stiffness matrix if none of the elements have mid-side nodes.

5.3   A Short Note on Automatic Data Generation

The reader will recall that the hypothetical user in Section 4 chose to read in his element coordinates and basic properties from data cards. This procedure can be time-consuming and expensive for problems involving large numbers of elements, especially when the element set geometry is not regular in the sense of Subsection 5.1 and Figure 13. However, it often happens that the element set geometry is regular in the sense that the element coordinates may be calculated from a general algorithm based upon the element string concept discussed in the previous subsection, i.e.:

```
DØ 20 LX=1, M
DØ 20 LY=1, N
LNUM=LY+N*(LX-1)
(X=F(LX, LY))
   .
   .
   .
```

where $\vec{X}=\{X1, Y1, X2, Y2,...\}$ is the element coordinate vector and $\vec{F}$ is a floating point function of LX and LY. An algorithm of this type may be used to generate the data base, as it is

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needed, for calculation of element stiffness matrices and element stresses. This results in the trade-off of a slight increase in execution time and some decrease in required storage space, since it is not necessary to carry large vectors of global nodal coordinates in core.

5.4 A Handy Trick

There often occur plane elasticity problems in which a large number of degrees of freedom are to have prescribed displacements. For example, suppose a rectangular domain such as the one shown in Figure 15 is to be analyzed, and that all four edges of the domain are clamped. This means that 25 to 30 per cent to the total degrees of freedom will have prescribed displacements.

A significant amount of execution time may be saved by modifying the numbering scheme shown in Figure 15, so that the edge degrees of freedom have the largest global numbers. Let NDT be the total number of degrees and NFT be the total number of unconstrained degrees. Then the modified number scheme assigns:
        1, 2,..., NFT
to the unconstrained degrees and:
        NFT+1, NFT+2,..., NDT
to the degrees along the edges of the structure. Program stages 1 through 5 are completed in standard fashion.

However, just before factoring **K** , the user may fool FEABL by inserting:
        ITEMP=NDT
        NDT=NFT
in his MAIN subprogram. Subroutines FACTPD/FACTSD and SIMULQ will then solve only for the unknown displacements 1, 2,..., NFT. At the beginning of Stage 7, when the full displacement vector may be required again, the user inserts:
        NDT=ITEMP

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The real value of this trick depends on a trade-off between execution time and core storage. The modified numbering scheme will result in a requirement for additional storage space for $\mathbf{K}$, over what is needed by the numbering scheme discussed in Subsection 5.2. The excess requirement is given by:

$$\Delta S \approx (NDT-NFT)\,(NDT-B)$$

where B is the average semi-bandwidth of the unconstrained part of $\mathbf{K}$.

5.5   Time-Saving Techniques for Design Studies

The "design study" approach to finite element analysis may take any of the following forms:

1. Consideration of a number of loading environments applied to a unique structure with unique displacement boundary conditions.

2. Analysis of a unique structure under various environments in which prescribed displacements are changed, as well as prescribed forces.

3. Consideration of variations in the structure itself to meet a given environment.

The first two categories are self-explanatory. An example of the third might be an analysis of a truss structure in which the compression bars are checked for buckling:

$$P_{Bar} < P_{cr} = \pi^2 E I / \ell^2$$

If any bar were found to be under a compressive load greater than its P_{cr}, the program might incorporate an algorithm for redesigning the bar and re-analyzing the new structure.

The above categories of problems can be studied using FEABL with only minor modifications. No changes are required in the data location interfaces and if a change to the process sequence interface is required, it involves only an intelligent

application of the rules presented above. However, the analyst
must be familiar with programming techniques required for com-
munication between his computer's core and external storage devices
(system disks, drums, tape units). These techniques vary from
installation to installation and will not be discussed in specific
terms here. The analyst should consult the operating manuals
applicable to the system which he will be using.

Load environment case study is the easiest type of multi-
solution problem to handle. Since the displacement boundary
conditions are not varied, only one assembly and factoring of
the master stiffness matrix need be done. The factored form
of K is held in the /DATA/ vector while each prescribed vector
is formed and the displacements are solved for. The quantity:

$$Q^* = \hat{Q}_E - K_{FD} \hat{u}_D$$

$\hat{Q}_E$ = Assembled Element Equivalent Nodal Forces
and the prescribed displacements $\hat{u}_D$ must be saved. These two
quantities are found in the Force/Displacement Vector:

$$Q = \left\{ \begin{matrix} Q^* \\ \hat{u}_D \end{matrix} \right\}$$

in the /DATA/ vector just after BCØN is called. (The prescribed
external loads Q_F, on the other hand, are not really needed
until SIMULQ is called.) Figure 16 illustrates the modified
FEABL process sequence which will accomplish this result.

The second category involves essentially the same techniques
as the first. However, in this case the data which must be saved
in external storage are:

1. $\hat{Q}_E$ – Assembled element equivalent nodal forces.
2. K – Assembled master stiffness matrix <u>prior to</u>
<u>boundary condition application.</u>

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FIGURE 16

For each case the user must input the contents of the constraint vector, fetch $\hat{Q}_E$, accumulate $\hat{Q}_F$ and reset $\hat{u}_D$, fetch $\mathbf{K}$ and finally re-enter the standard FEABL process at the point where BCØN is called.

Significant time savings are gained in the third problem category if the design changes involve only a few elements in a structure composed of a large number of elements. Suppose that the displacement solution has been found from the force-displacement relations for the initial structure:

$$ \mathbf{K}u = \hat{Q}_F + Q^\star $$

During calculation of the element stresses, it is discovered that design changes are required in one or more elements; these changes will appear in the force-displacement relations as a modification $\Delta \mathbf{K}$ of the master stiffness matrix and (possibly) a modification $\Delta \hat{Q}_E$ of the assembled element nodal force vector.

The new force-displacement relations:

$$\left(\mathbf{K} + \Delta\mathbf{K}\right) u' = \hat{Q}_F + Q^* + \Delta Q^*$$

must then be solved for the modified displacements u'.

To program the design change technique with FEABL, the user must take the following nonstandard actions:

1. Output $\mathbf{K}$ and $\hat{Q}_F + \hat{Q}_E$ to external storage after assembly and input of prescribed quantities, but before calling BCØN.

2. After the initial displacement solution u has been obtained, transfer it to temporary core storage outside the /DATA/ vector. (The Force/Displacement Vector area in the /DATA/ vector will be required for accumulation of $\Delta\hat{Q}_E$.) Zero the Force/Displacement Vector and the Master Stiffness Matrix Array.

3. Calculate stresses from u, element-by-element. (The user will need his own version of XTRACT to extract the proper u_{el} from temporary core storage.) When a design change is required, calculate $\Delta\mathbf{k}_{el}$ and $\Delta\hat{Q}_{el}$ and accumulate them using FEABL subroutine ASEMBL.

4. Apply rotation transformations (if any) to $\Delta\mathbf{K}$ and $\Delta\hat{Q}_E$, reset entries of $\Delta\hat{Q}_E$ to zero where displacements are prescribed.

5. Fetch $\mathbf{K}$ and $\hat{Q}_F + \hat{Q}_E$ from external storage and accumulate them to $\Delta\mathbf{K}$ and $\Delta\hat{Q}_E$. Apply boundary conditions and solve for u'.

The above process may be repeated in an iterative design process, with $\mathbf{K} + \Delta\mathbf{K}$, $\hat{Q}_F + Q^* + \Delta Q^*$ at the beginning of each new design step playing the role of the initial values $\mathbf{K}$, $\hat{Q}_F + Q^*$. This procedure is referred to as iterative updating, and is also applicable to nonlinear elastic and elastic-plastic analysis of continua. In these types of analysis the "design change"

results from following a material stress-strain curve and/or
testing the satisfaction of a yield conditon (e.g., the Mises-
Hencky criterion).

5.6   Substructuring with FEABL

The FEABL software system has been designed primarily
for in-core solutions. Up to 1,500 degrees of freedom can be
handled on currently available hardware with 500 KBYTE (125 KWord)
memory. FEABL's in-core capability may be extended by means
of the substructuring technique outlined briefly here.

Figure 17 illustrates a domain which has been divided
into a small number of substructures. The desired stress solution
accuracy is on a scale much smaller than the substructure dimension,

Substructure divided into ordinary elements

FIGURE 17

so each substructure is subdivided further into ordinary elements.
Let the subscripts I and B refer, respectively, to the interior
and boundary degrees of freedom in a substructure. Then the

force-displacement relations for the substructure may be partitioned into:

$$\begin{bmatrix} K_{BB} & K_{BI} \\ K_{BI}^T & K_{II} \end{bmatrix} \begin{Bmatrix} u_B \\ u_I \end{Bmatrix} = \begin{Bmatrix} Q_B \\ Q_I \end{Bmatrix}$$

The interior degrees of freedom may be eliminated by the process of static condensation, which transforms the substructure force-displacement relations to the form:

$$K_{BB}^c \, u_B = \left(K_{BB} - K_{BI} K_{II}^{-1} K_{BI}^T \right) u_B = Q_B - K_{BI} K_{II}^{-1} Q_I = Q_B^c$$

Each substructure may now be considered as a "superelement" having degrees of freedom only on the interelement boundaries, as shown in Figure 18. The quantities K_{BB}^c and Q_B^c are then

FIGURE 18

assembled in the same manner as the stiffness matrices and equivalent nodal force vectors for ordinary elements. After the displacement solution on the substructure boundaries has been obtained, the original force-displacement relations for each substructure may be used, with u_B as prescribed displacements, to obtain the interior solution.

In addition to extension of FEABL's problem size capability, the substructuring technique may also be used to reduce roundoff error and to save execution time. All of these improvements depend upon optimization of the relative numbers of substructures and of ordinary elements within the substructures. Use of the substructuring technique with FEABL software will be discussed in detail in a future publication.

5.7   Error Estimation Methods

Rigorous mathematical proofs exist showing that a finite element analysis for the stress and displacement distributions in a structure converges to the exact solution as the number of elements in a given domain is increased, provided only that certain easily satisfied restrictions are obeyed (Ref.3). However, since all calculations done in a digital computer are imprecise, errors due to roundoff will occur. It has been shown (Ref. 4) that roundoff error increases in an extremely complex way as the number of elements is increased. Therefore, it is advisable to resort to some numerical method which will produce a reasonable estimate of the error, and which does not depend upon the details of what types of elements are used or what boundary conditions are applied to the structure. Four methods are presented here.

5.7.1 Irons' Energy Variance Criterion

The energy criterion proposed by Irons (Ref. 5) is the most economical, requiring no more time than the calculation of the rounding error parameter discussed in Subsection 3.2.6. The diagonal entries K_{ii} of the master stiffness matrix must be saved, either in temporary core storage or on an external unit. After the displacement solution has been obtained, the energy variance can be calculated from:

$$\mathscr{C} = 2^{-p+1} \sqrt{B} \frac{\sum\limits_{i=1}^{NDT} K_{ii} (u_i)^2}{\sum\limits_{i=1}^{NDT} \sum\limits_{j=1}^{NDT} K_{ij} u_i u_j}$$

where p is the computer precision in decimal places and B is
the average semi-bandwidth of K . The value of B can be obtained
easily from the address index parameters*:

B=FLØAT(LK+1-IK)/FLØAT(NDT)

One half the value of the denominator of the energy variance
expression is provided to the user by FEABL subroutine SIMULQ
in the argument ENERGY of that subroutine.

Some care must be exercised in applying Irons' energy
criterion. The user will obtain unrealistically large values
of $\mathscr{C}$ in problems in which the structure is constrained very
lightly, and in which large amounts of "self-energy" can be stored
when a single degree of freedom is displaced. The cantilever
beam is a good example of a structure to which the Irons' criterion
cannot be applied.

5.7.2 The Residual Force Method

Calculation of residual forces provides a more detailed
picture of the distribution of errors through the structure.
The master stiffness matrix may be saved in external storage for
this purpose immediately after calling BCØN. Approximate values
of the reaction forces at degrees where displacements were pre-
scribed may be obtained as well by saving K just prior to calling
BCØN.

---

*FLØAT is an IBM library function which converts integers to
floating point numbers.

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Let u^*_F be the approximate displacement solution obtained from the constrained force-displacement relations:

$$\begin{bmatrix} K_{FF} & 0 \\ 0 & I \end{bmatrix} \begin{Bmatrix} u_F \\ \hat{u}_D \end{Bmatrix} = \begin{Bmatrix} \hat{Q}_F - K_{FD}\hat{u}_D \\ \hat{u}_D \end{Bmatrix}$$

Then:

$$Q^* = \begin{Bmatrix} Q^*_F \\ Q^*_D \end{Bmatrix} = K \begin{Bmatrix} u^*_F \\ \hat{u}_D \end{Bmatrix}$$

is the approximate force vector. When K has been returned to core after the displacement solution has been obtained, Q^* can be calculated by the following algorithm:

```
      DIMENSIØN R(500)
C     VECTØR R MUST ALLØW ENØUGH STØRAGE FØR THE FULL
C     FØRCE VECTØR Q-STAR IF REACTIØN FØRCES ARE
C     DESIRED.  ØTHERWISE, ENØUGH FØR THE UNCØNSTRAINED
C     DEGREES ØF FREEDØM MUST BE ALLØWED
C     ALLOCATE EXT:RNAL FILES
      :
      :
      (Standard FEABL process sequence, except where noted)
      :
      :
      :
C     K MATRIX TØ EXTERNAL STØRAGE
      WRITE (...) (REAL(I), I=IK,LK)
C     FØRCE/DISPL MUST ALSØ BE SAVED FØR CØMPARISØN LATER
      WRITE (...) (REAL(I), I=IQ, LQ)
      CALL BCØN
      CALL FACTPD
      CALL SIMULQ(STRE)
      :
      :
      (Stress solution, if desired)
      :
      :
C     RETURN K MATRIX FØR CALCULATIØN ØF Q-STAR
      READ (...) (REAL(I), I=IK, LK)
C     INITIALIZE CØNSTRAINED RØW PØINTER AT 1ST NØNZERØ RØW
      DØ 50 II=ICØN, LCØN
      IF(INTGR(II) .EQ. 0) GØ TØ 50
      NC=II
      GØ TØ 60
50    CØNTINUE
```

```
C     INITIALIZE R VECTØR PØINTER
   60 NR=1
C     LØØP ØVER DEGREES ØF FREEDØM, SKIPPING CØNSTRAINED DEGREES
      DØ 70 IRØW=1, NDT
      IF(INTGR(NC) .NE. IRØW) GØ TØ 71
C     UPDATE RØW PØINTER
      NC=NC+1
      GØ TØ 70
C     FØRM SUMMATIØN ØF K(IRØW, J)* U(J), J=LNZ TØ NDT
   71 INIT=ILNZ+IRØW-1
      INIT=INTGR(INIT)
      KK=IKØUNT+IRØW-1
      KK=INTGR(KK)
      SUM=0.
      DØ 72 J=INIT, IRØW
      KADR=KK+J
      JJ=IQ+J-1
   72 SUM=SUM+REAL(KADR)* REAL(JJ)
C     REMAINDER ØF SUM MUST TAKE K ENTRIES FRØM CØL IRØW
      INIT=IROW+1
      IF(INIT .GT. NDT) GØ TØ 74
      DØ 73 J=INIT, NDT
C     MAKE SURE RØW J HAS AN ENTRY IN CØL IRØW
      JJ=ILNZ+J-1
      IF(INTGR(JJ) .GT. IRØW) GØ TØ 73
      KADR=IKØUNT+J-1
      KADR=INTGR(KADR)+IRØW
      JJ=IQ+J-1
      SUM=SUM+REAL(KADR)*REAL(JJ)
   73 CØNTINUE
C     PLACE SUM IN NEXT AVAIL R LØCATIØN AND UPDATE R PØINTER
   74 R(NR)=SUM
      NR=NR+1
   70 CØNTINUE
C     RETURN ØRIGINAL Q TØ CØRE FØR CØMPARISØN
      READ (...) (REAL(I), I=IQ, LQ)
C     CALCULATE RESIDUAL FØRCES Q-(Q-STAR)=Q-R
C     REINITIALIZE NC AND NR
      NC=II
      NR=1
      DØ 80 IRØW=1, NDT
      IF(INTGR(NC) .NE. IRØW) GØ TØ 81
      NC=NC+1
      GØ TØ 80
   81 JJ=IQ+IRØW-1
      R(NR)=REAL(JJ)-R(NR)
      WRITE (KW, 500) IRØW, R(NR)
  500 FØRMAT (20H RESIDUAL FØRCE NØ.,I6, 1X, 1H=, E10.3)
      NR=NR+1
   80 CØNTINUE
```

In the above algorithm, the vector of residual forces:

$$\mathbf{R}_F = \hat{Q}_F - Q_F^* = \Delta Q_F$$

has been calculated. The entries of $\mathbf{R}_F$ provide the detailed picture mentioned at the beginning of this subsection. Another useful parameter for overall error measurement is the force vector magnitude ratio:

$$r_F = \left[\frac{\Sigma (R_i^2)}{\Sigma (\hat{Q}_i - \Sigma K_{ij} u_j)^2} \right]^{1/2} = \frac{|\Delta Q_F|}{|\hat{Q}_F|}$$

where the summations extend only over the unconstrained degrees of freedom.

5.7.3 Re-Solution for Residual Displacements

Although the residual force vector is fairly easy to calculate, the interpretation of its meaning is not a trivial task. The displacement solution must certainly be judged acceptable if, for example, there are many residual forces on the order of 1 lb. at degrees where forces of 1,000 lb. were applied originally. However, a common situation in finite element analysis is that the applied force is zero at many degrees of freedom. What does a 1 lb. residual force mean at these points? The averaged measure r_F presented in the previous section relieves this detailed interpretation problem to some degree. However, the averaged measure of greatest interest to the analyst is the displacement vector magnitude ratio:

$$d_F = \frac{|\Delta u_F|}{|u_F|}$$

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Unfortunately, no simple relation exists between d_F and the force vector magnitude ratio r_F. If $\mathbf{K}$ is an ill-conditioned matrix, d_F may in fact be much larger than r_F.

If the analyst is willing to spend some additional computing time, the displacement residuals and an approximate calculation of d_F may be obtained by re-solution. This technique requires an additional external storage file capable of holding $\mathbf{K}$. Just before $\mathbf{K}$ is returned to core in the algorithm of Subsection 5.7.2, its factored form $\mathbf{LDL}^T$ is read into this extra file. Then, picking up where the previous algorithm ended, the original displacement vector magnitude is calculated, $\mathbf{LDL}^T$ is returned to core and the contents of $\mathbf{R}$ are transferred to the proper locations in the force/displacement block of the /DATA/ vector. Re-solution is now done simply by calling SIMULQ again, after which $\Delta \boldsymbol{u}_F$ will be found in the force/displacement block. The displacement vector magnitude ratio obtained from this procedure is actually:

$$ d_F^* = \frac{|\Delta \boldsymbol{u}_F|}{|\boldsymbol{u}_F^*|} $$

and $d_F^* \neq d_F$ unless $|\Delta \boldsymbol{u}_F| << |\boldsymbol{u}_F^*|$.

5.7.4 The Method of Rigid Body Modes

A somewhat less cumbersome technique for error measurement can be employed when a structure is modeled by elements which contain a full set of rigid body modes in their assumed displacement fields. Let

$$ \boldsymbol{r} = \left\{ \begin{matrix} \boldsymbol{r}_F \\ \hat{\boldsymbol{r}}_D \end{matrix} \right\} $$

be any rigid body displacement vector for the whole structure (e.g., unit vertical translation at every node). Then if $\boldsymbol{r}$ is introduced into the unconstrained force-displacement relations and the right hand side is calculated, there will result:

$$\mathbf{K}\,\boldsymbol{r} = \begin{Bmatrix} 0 \\ 0 \\ \vdots \\ \vdots \\ 0 \end{Bmatrix}$$

to the accuracy of the user's computer. Conversely, solution
of the constrained force-displacement relations:

$$\begin{bmatrix} \mathbf{K}_{FF} & 0 \\ 0 & \mathbf{I} \end{bmatrix} \begin{Bmatrix} u_F \\ \hat{r}_D \end{Bmatrix} = \begin{Bmatrix} -\mathbf{K}_{FD}\,\hat{r}_D \\ \hat{r}_D \end{Bmatrix}$$

would result in $\mathcal{U}_F = \boldsymbol{r}_F$ to the accuracy of the computer if
there were no roundoff error. Now, since the exact $\boldsymbol{r}_F$ can
be inferred from $\hat{\boldsymbol{r}}_D$ merely by inspection, the above problem
may be solved as an auxiliary to the real problem, and the error
measure:

$$d_{RB} = \frac{|\boldsymbol{r}_F - \mathcal{U}_F|}{|\boldsymbol{r}_F|}$$

may be calculated. The rigid body mode technique requires the
saving only of $\mathbf{K}$ (before constraint). A simple algorithm will
serve to calculate $\boldsymbol{r}_F$ and $\boldsymbol{r}_F - \mathcal{U}_F$, and the additional execution
time required amounts only to calling each of BCØN, FACTPD (or
FACTSD) and SIMULQ once extra. Also, the lengthy residual force
and re-solution algorithm in the MAIN program is avoided.
Recent tests of the method of rigid body modes on a cantilever
beam have shown that when $d_F < 0.1$, $d_{RB} << d_F$. However, when $d_F > 0.1$
(the region of primary interest) d_{RB} performs as well as d_F^*.

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APPENDIX A

CØMMØN AREAS DATA REQUIREMENTS

The following table summarizes what information is expected by each FEABL subroutine in the four control parameter CØMMØN areas and in the /DATA/ vector CØMMØN area.

CØMMØN AREA \ SUBROUTINE NAME	ASEMBL	BCØN	FACTPD/FACTSD	ØRK	RØTATE	SETUP	SIMULQ	XTRACT
/IO/ (Printer Code KW Only)	✓	✓	✓	✓	✓	✓	✓	
/SIZE/ NET, NDT		✓	✓	✓	✓	✓	✓	
/BEGIN/ Address index parameters	✓	✓	✓	✓	✓		✓	✓
/END/ Address index parameters		✓	✓	✓ Except LK			✓	
/DATA/ Vector:								
1. Constraint Vector		✓	✓			(c)	✓	
2. Address Count Vector	✓	✓	✓		✓		✓	
3. LNZ Vector	✓	✓	✓		✓		✓	
4. Assembly List	✓			✓				✓
5. F/D Vector	(a)	✓			Only $\hat{Q}_E$ (b)		✓	u
6. K Matrix	(a)	✓	✓		✓		LDL^T	

(a) These blocks are zeroed by ØRK before the first element is assembled.

(b) $\hat{Q}_E$ = Vector of assembled element equivalent nodal forces.

(c) Block zeroed by SETUP

APPENDIX B

APPLYING SHOEHORN AND STOPWATCH

This appendix contains data from which the user may estimate the total amount of core storage required by a FEABL-based program and the approximate execution time (CPU time) the run will take. Such estimates will prove useful aids in making trade-off decisions. The data given in this appendix is based on runs done on an IBM 370/155 using the FØRTPAN G compiler. The numbers will vary somewhat from one machine or compiler to another.

B.1   Estimation of Core Storage Requirement

The following table gives the length of each FEABL subroutine in BYTES (as compiled in FØRTRAN G) and words, and the deck size. On IBM 360 and 370 series hardware, core storage calculations are normally done in terms of BYTES, while words are used on many

Subroutine Name	Words	BYTES	No. of Cards in Deck*
ASEMBL	410	1,642	58
BCØN	599	2,396	125
FACTPD/FACTSD	669	2,674	126
ØRK	517	2,070	106
RØTATE	1,858	7,430	316
SETUP	726	2,904	94
SIMULQ	600	2,398	127
XTRACT	130	522	23
FEABL Software-Total	5,509	22,036	975

other hardware systems. The total storage requirement for programs plus data can be estimated as follows:

---

*Includes all comment cards

Item:	KWords	KBYTES
1. FEABL Software	5.5	22.1
2. User programs (a)	2.5	10.0
3. System library subprograms (b)	5.1	22.0
4. /DATA/ Vector (c)	(L)	(4L)
5. Four control parameter /CØMMØN/ areas	--	0.1
Totals	13.1+(L)	53.2+(4L)

(a)   Estimate for a typical analysis with a MAIN program and two generator subroutines.

(b)   Includes library functions such as SIN, CØS, SQRT and systems management subroutines.

(c)   L=the dimension of the /DATA/ vector in KWords (1000 words).

B.2   Estimation of CPU Time Requirement

The process of factoring the master stiffness matrix into its triple product:

$$K = LDL^T$$

is the primary time consumer in any finite element analysis. Some study of the algorithms in FEABL subroutine FACTPD/FACTSD will convince the reader that the CPU time consumed is proportional to NB^2, where N is the total number of unconstrained DOF in the assembled structure and B is the average semi-bandwidth of the master stiffness matrix. B may be calculated approximately from N and the population density of K :

$$B \doteq \frac{(N+1)P}{2}$$

where

$$P=\text{Population Density} = \frac{\text{Total No. of Stored Entries}}{\text{Total Entries in Full Lower Triangle}}$$

Experience on the IBM 370/155 indicates that the required CPU time is given approximately by:

$$\text{time} \doteq NB^2 \times 10^{-6} \text{ minutes}$$

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APPENDIX C

FØRTRAN IV LISTING OF FEABL SOFTWARE

The eight subroutines of the FEABL software system are listed in alphabetical order in this appendix:

Subroutine	Page
ASEMBL	71
BCØN	73
FACTPD/FACTSD	77
ØRK	81
RØTATE	84
SETUP	93
SIMULQ	96
XTRACT	100

The code is for FEABL Version 1 Release 1, with a 10,000-word /DATA/ vector. The following actions will convert the program to Version 2 ("On-Line") as explained in Section 2.3.

SUBROUTINE	Change DIMENSIØNs of REAL, INTGR to 2	Delete Declaration
ASEMBL	Card No. 0007	Card Nos. 0017 and 0020
BCØN	0007	0012 and 0015
FACTPD/FACTSD	0008	0013 and 0016
ØRK	0015	0020 and 0023
RØTATE	0007	0013 and 0017
SETUP	0007	0013 and 0015
SIMULQ	0007	0012 and 0015
XTRACT	0007	0011 and 0013

```
      SUBROUTINE ASEMBL(LNUM,NDE,ELK,ELQ)                              ASEM0001
C**********************************************************************ASEM0002
C FINITE ELEMENT ANALYSIS BASIC LIBRARY SUBROUTINE                    ASEM0003
C AEROELASTIC AND STRUCTURES RESEARCH LABORATORY                      ASEM0004
C MASSACHUSETTS INSTITUTE OF TECHNOLOGY                               ASEM0005
C**********************************************************************ASEM0006
      DIMENSION REAL(10000),INTGR(10000)                              ASEM0007
      DIMENSION ELK(NDE,NDE), ELQ(NDE)                                ASEM0008
C**************************                                            ASEM0009
C CHANGE THIS DIMENSION STATEMENT IF YOU HAVE RECEIVED AN ERROR MESSAGEASEM0010
C INDICATING LENGTH OF MNUM VECTOR HAS BEEN EXCEEDED.  ALSO CHANGE    ASEM0011
C INSTRUCTION INDICATED BELOW                                         ASEM0012
      DIMENSION MNUM(100)                                             ASEM0013
C**************************                                            ASEM0014
      COMMON /IO/ KR, KW, KP, KT1, KT2, KT3                           ASEM0015
      COMMON /BEGIN/ ICON,IKOUNT,ILNZ,IMASTR,IQ,IK                    ASEM0016
      COMMON /DATA/ REAL                                              ASEM0017
C VERSION 1 RELEASE 1 AUGUST 1972                                     ASEM0018
C                                                                     ASEM0019
      EQUIVALENCE (REAL(1), INTGR(1))                                 ASEM0020
C PRINT CONTROL                                                       ASEM0021
  901 FORMAT(1H0,I10,27H DOF/ELEMENT EXCEEDS LENGTH,I4,26H SPECIFIED FORASEM0022
     1 MNUM VECTOR,/,1X,44HERROR OCCURED DURING ASSEMBLY OF ELEMENT NO.,ASEM0023
     2 I10,/,1X,95HUSER MUST CORRECT CALL ARGUMENTS IN MAIN OR CHANGE LEASEM0024
     3NGTH OF MNUM VECTOR IN SUBROUTINE ASEMBL,/,1X,20HEXECUTION TERMINAASEM0025
     4TED)                                                            ASEM0026
      IKOUM1 = IKOUNT-1                                               ASEM0027
      IQM1 = IQ-1                                                     ASEM0028
C**************************                                            ASEM0029
C VALUE OF LENGTH MUST = DIMENSION OF MNUM                            ASEM0030
      LENGTH = 100                                                    ASEM0031
C**************************                                            ASEM0032
      IF (NDE .LE. LENGTH) GO TO 1                                    ASEM0033
      WRITE (KW,901) NDE, LENGTH, LNUM                                ASEM0034
      STOP                                                            ASEM0035
C GET MASTER ROW/COL NOS FOR ELEMENT AND STORE IN MNUM               ASEM0036
```

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```
  1 INDEX = IMASTR+LNUM-1                                              ASEM0037
    DO 2 I = 1,NDE                                                     ASEM0038
    J = INTGR(INDEX)+I-1                                               ASEM0039
  2 MNUM(I) = INTGR(J)                                                 ASEM0040
C LOOP OVER ROWS OF ELO AND OVER LOWER TRIANGLE OF ELEMENT K MATRIX    ASEM0041
    DO 4 LROW = 1,NDE                                                  ASEM0042
    INDEX = IQM1+MNUM(LROW)                                            ASEM0043
C ASSEMBLE ELEMENT EQUIVALENT NODAL FORCE INTO Q VECTOR                ASEM0044
    REAL(INDEX) = REAL(INDEX)+ELQ(LROW)                                ASEM0045
    DO 4 LCOL = 1,LROW                                                 ASEM0046
    MROW = MNUM(LROW)                                                  ASEM0047
    MCOL = MNUM(LCOL)                                                  ASEM0048
    IF (MROW .GE. MCOL) GO TO 3                                        ASEM0049
    MROW = MNUM(LCOL)                                                  ASEM0050
    MCOL = MNUM(LROW)                                                  ASEM0051
C CALCULATE ABSOLUTE ADDRESS OF K(MROW,MCOL)                           ASEM0052
  3 INDEX = IKJUM1+MROW                                                ASEM0053
    KADR = INTGR(INDEX)+MCOL                                           ASEM0054
C ASSEMBLE STIFFNESS COEFFICIENT                                       ASEM0055
  4 REAL(KADR) = REAL(KADR)+ELK(LROW,LCOL)                             ASEM0056
    RETURN                                                             ASEM0057
    END                                                                ASEM0058
```

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```
      SUBROUTINE BCON                                              BCON0001
C*****************************************************************BCON0002
C FINITE ELEMENT ANALYSIS BASIC LIBRARY SUBROUTINE               BCON0003
C AEROELASTIC AND STRUCTURES RESEARCH LABORATORY                 BCON0004
C MASSACHUSETTS INSTITUTE OF TECHNOLOGY                          BCON0005
C*****************************************************************BCON0006
      DIMENSION REAL(10000),INTGR(10000)                          BCON0007
      COMMON /IO/ KR, KW, KF, KT1, KT2, KT3                        BCON0008
      COMMON /SIZE/ NET, NOT                                       BCON0009
      COMMON /BEGIN/ ICON,IKOUNT,ILNZ,IMASTR,IQ,IK                 BCON0010
      COMMON /END/ LCON,LKOUNT,LLNZ,LMASTR,LQ,LK                   BCON0011
      COMMON /DATA/ REAL                                           BCON0012
C                                                                 BCON0013
C VERSION 1 RELEASE 1 AUGUST 1972                                 BCON0014
C                                                                 BCON0015
      EQUIVALENCE (REAL(1), INTGR(1))                             BCON0016
C PRINT CONTROL                                                   BCON0017
  901 FORMAT(79HODISPLACEMENT CONSTRAINTS HAVE BEEN APPLIED TO THE FOLLOBCON0018
     1WING DEGREES OF FREEDOM,/,5X,7HDOF NO.,2X,12HDISPLACEMENT)   BCON0019
  902 FORMAT (2X,I10,2X,E10.3)                                     BCON0020
  903 FORMAT(1X,72H***************************A8OVE DOF NUMBER APPEARS TWICEBCON0021
     1 IN CONSTRAINT LIST,/,1X,85HEXECUTICN TERMINATED IN SUBROUTINE BCOBCON0022
     2N DUE TO POSSIBLITY OF BOUNCARY CONDITION ERROR)             BCON0023
  904 FORMAT(62HOYOUR STRUCTURE IS FLYING FREE. PLEASE CONSTRAIN IT NEXTBCON0024
     1 TIME.,/,1X,28HEXECUTION TERMINATED IN BCON)                 BCON0025
      IKOUM1 = IKOUNT-1                                            BCON0026
      ILNZM1 = ILNZ-1                                              BCON0027
      IQM1 = IQ-1                                                  BCON0028
C PRINT ENTRY MESSAGE                                             BCON0029
      WRITE (KW,901)                                               BCON0030
C ORDER THE CONSTRAINT ROW NUMBERS IN ASCENDING SEQUENCE          BCON0031
      LAST = LCON-1                                                BCON0032
      IF (LCON .GT. 0) GO TO 1                                     BCON0033
      WRITE (KW,904)                                               BCON0034
      STOP                                                         BCON0035
    1 IFLAG = 0                                                    BCON0036
      DO 2 I = ICON,LAST
```

```
      IF (INTGR(I+1) .GE. INTGR(I)) GO TO 2                              BCON0037
      J = INTGR(I)                                                      BCON0038
      INTGR(I) = INTGR(I+1)                                             BCON0039
      INTGR(I+1) = J                                                    BCON0040
      IFLAG = 1                                                         BCON0041
    2 CONTINUE                                                          BCON0042
      IF (IFLAG .EQ. 1) GO TO 1                                         BCON0043
C CHECK TO SEE IF ANY ROW NUMBERS HAVE BEEN ENTERED IN CONSTRAINT       BCON0044
C VECTOR - ABORT THE RUN IF NONE HAVE BEEN                              BCON0045
      J = 0                                                             BCON0046
      DO 100 I = ICON,LCON                                              BCON0047
  100 J = J+INTGR(I)                                                    BCON0048
      IF (J .GT. 0) GO TO 200                                           BCON0049
      WRITE (KW,904)                                                    BCON0050
      STOP                                                              BCON0051
C OUTPUT CONSTRAINT LIST                                                BCON0052
  200 DO 4 I = ICON,LCON                                                BCON0053
      IF (INTGR(I) .EQ. 0) GO TO 4                                      BCON0054
C CHECK FOR REPEATED DOF AFTER 1ST ONE                                  BCON0055
      IF (I .EQ. ICON) GO TO 3                                          BCON0056
      IF (INTGR(I) .NE. INTGR(I-1)) GO TO 3                             BCON0057
      WRITE (KW,903)                                                    BCON0058
      STOP                                                              BCON0059
    3 J = IQM1+INTGR(I)                                                 BCON0060
      WRITE (KW,902) INTGR(I), REAL(J)                                  BCON0061
    4 CONTINUE                                                          BCON0062
C LOOP OVER CONSTRAINED DOF FOR MODIFICATION OF COMPLETELY ASSEMBLED    BCON0063
C FORCE VECTOR                                                          BCON0064
      DO 71 I = ICON,LCON                                               BCON0065
      IF (INTGR(I) .EQ. 0) GO TO 71                                     BCON0066
C CHECK IF PRESCRIBED DISPLACEMENT = 0 -- IF IT DOES, SKIP FORCE VECTOR BCON0067
      MROW = INTGR(I)                                                   BCON0068
      M = IQM1+MROW                                                     BCON0069
      IF (REAL(M) .EQ. 0.) GO TO 71                                     BCON0070
C DISPL .NE. 0 -- LOOP OVER ALL ROWS TO MODIFY FORCE VECTOR -- SKIP     BCON0071
C CONSTRAINED ROWS (CONTROLLED BY VALUE OF NEXT)                        BCON0072
```

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```
          NEXT = ICON
          DO 7 NROW = 1,NDT
          IF (NROW .NE. INTGR(NEXT)) GO TO 5
          NEXT = NEXT+1
          GO TO 7
        5 IF (MROW .GT. NROW) GO TO 51
C CHECK FOR COUPLING OF ROW NROW WITH COL MROW
          J = ILN2M1+NROW
          IF (INTGR(J) .GT. MROW) GO TO 7
          IROW = NROW
          ICOL = MROW
          GO TO 6
C CHECK FOR COUPLING OF ROW MROW WITH COL NROW
       51 J = ILN2M1+MROW
          IF (INTGR(J) .GT. NROW) GO TO 7
          IROW = MROW
          ICOL = NROW
C SUBTRACT K*(PRESCR DISPL) FROM FORCE VECTOR
        6 KADR = IKOUM1+IROW
          KADR = INTGR(KADR)+ICOL
          N = IOM1+NROW
          REAL(N) = REAL(N)-REAL(KADR)*REAL(M)
        7 CONTINUE
       71 CONTINUE
C LOOP OVER CONSTRAINED ROWS TO DECOUPLE THEM FROM REST OF K MATRIX
          DO 11 I = ICON,LCON
          IF (INTGK(I) .EQ. 0) GO TO 11
          MROW = INTGK(I)
          INIT = ILN2M1+MROW
          INIT = INTGR(INIT)
C SET ROW = 0
          M = IKOUM1+MROW
          M = INTGR(M)
          DO 8 MCOL = INIT,MROW
          KADR = M+MCOL
        8 REAL(KADR) = 0.
```

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```
C SET COLUMN = 0 IN ROWS WHOSE LNZE COL NC IS .LE. MROW -- SKIP THIS     BCON0109
C SECTION IF MROW .EQ. THE LAST ROW                                      BCON0110
      IF (MROW .EQ. NDT) GO TO 10                                        BCON0111
      INIT = MROW+1                                                      BCON0112
      DO 9 NROW = INIT,NDT                                               BCONO113
      N = ILNZAL+NROW                                                    BCON0114
      IF (INTGR(N) .GT. MROW) GO TO 9                                    BCON0115
      KADR = IKUUM1+NROW                                                 BCON0116
      KADR = INTGR(KADR)+MROW                                            BCON0117
      REAL(KADR) = 0.                                                    BCON0118
    9 CONTINUE                                                           BCON0119
C SET DIAGONAL ENTRY = 1                                                 BCON0120
   10 KADR = M+MROW                                                      BCON0121
      REAL(KADR) = 1.                                                    BCON0122
   11 CONTINUE                                                           BCON0123
      RETURN                                                             BCON0124
      END                                                                BCON0125
```

```
          SUBROUTINE FACTSD                                             FACT0001
C***********************************************************************FACT0002
C FINITE ELEMENT ANALYSIS BASIC LIBRARY SUBROUTINE                      FACT0003
C AEROELASTIC AND STRUCTURES RESEARCH LABORATORY                        FACT0004
C MASSACHUSETTS INSTITUTE OF TECHNOLOGY                                 FACT0005
C***********************************************************************FACT0006
C THIS IS ALSO SUBROUTINE FACTPD - SEE ENTRY POINT BELOW                FACT0007
      DIMENSION REAL(10000), INTGR(10000)                               FACT0008
      COMMON /IO/ KR, KW, KP, KIL, KT2, KT3                             FACT0009
      COMMON /SIZE/ NET, NDT                                            FACT0010
      COMMON /BEGIN/ ICON,IKOUNT,ILNZ,IMASTR,IQ,IK                      FACT0011
      COMMON /END/ LCON,LKOUNT,LLNZ,LMASTR,LQ,LK                        FACT0012
      COMMON /DATA/ REAL                                                FACT0013
C VERSION 1 RELEASE 1 AUGUST 1972                                       FACT0014
C                                                                       FACT0015
      EQUIVALENCE (REAL(1), INTGR(1))                                   FACT0016
C PRINT CONTROL                                                         FACT0017
  901 FORMAT (1H0,47X,25HTRIPLE FACTOR ENTRY POINT,/,53H K MATRIX NOT POFACT0018
     2SITIVE-DEFINITE IN THE FOLLOWING ROWS )                           FACT0019
  902 FORMAT (40X,I12)                                                  FACT0020
  903 FORMAT (48X,4HNONE )                                              FACT0021
  904 FORMAT (58H USER SPECIFIED SEMI-DEFINITE MATRIX.  EXECUTION CONTINFACT0022
     2UES )                                                             FACT0023
  905 FORMAT (40H USER SPECIFIED POSITIVE-DEFINITE MATRIX,/,21H EXECUTIOFACT0024
     2N TERMINATES )                                                    FACT0025
  906 FORMAT (25H0K MATRIX SINGULAR IN ROW,I12,/,21H EXECUTION TERMINATEFACT0026
     2S )                                                               FACT0027
  907 FORMAT (61H0LARGEST ROUNDING ERROR IN DIAGONAL FACTORING OCCURRED FACT0028
     2IN ROW,I12,/,43H NUMBER OF LOWEST SIGNIFICANT FIGURES LOST=,I3)    FACT0029
  908 FORMAT (42H ROUNDING ERROR EXCEEDS ACCEPTABLE MAXIMUM,/,21H EXECUTFACT0030
     2ION TERMINATES )                                                  FACT0031
C SET ENTRY FLAG                                                        FACT0032
      NPD = 0                                                           FACT0033
      GO TO 1                                                           FACT0034
      ENTRY FACTPD                                                      FACT0035
      NPD = 1                                                           FACT0036
```

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```
C PRINT ENTRY MESSAGE
1 WRITE (KW,9011)
  IKOUM1 = IKOUNT-1
  ILNZM1 = ILNZ-1
C INITIALIZE AT 1ST NONZERO ENTRY IN CONSTRAINT VECTOR
  DO 2 I = ICON,LCON
  IF (INTGR(I) .EQ. 0) GO TO 2
  GO TO 3
2 CONTINUE
3 NEXT = 1
  IF (INTGR(NEXT) .EQ. 1) NEXT = NEXT+1
C INITIALIZE ERROR PARAMETERS
  JROW = 1
  MPD = 0
  TEST = 1.
C DO FIRST ROW AS SPECIAL CASE
  IF (REAL(IK) .EQ. 0.) GO TO 14
  IF (REAL(IK) .GT. 0.) GO TO 4
C NPD MESSAGE AND FLAG
  WRITE (KW,9002) JROW
  MPD = 1
C LOOP OVER REMAINING ROWS
4 DO 13 MROW = 2,NDT
C CHECK FOR CONSTRAINED ROW - SKIP IF FOUND AND RESET FOR THE NEXT ONE
  IF (MROW .NE. INTGR(NEXT)) GO TO 5
  NEXT = NEXT+1
  GO TO 13
C FREE ROW - FACTOR FROM LNZ COL NO TO ROW NO
5 M = ILNZM1+MROW
  M = INTGR(M)
  MM = IKOJM1+MROW
  MM = INTGR(MM)
  DO 12 MCOL = M,MROW
  SUM = 0.
  NN = IKOJM1+MCOL
  NN = INTGR(NN)
```

```
FACT0037
FACT0038
FACT0039
FACT0040
FACT0041
FACT0042
FACT0043
FACT0044
FACT0045
FACT0046
FACT0047
FACT0048
FACT0049
FACT0050
FACT0051
FACT0052
FACT0053
FACT0054
FACT0055
FACT0056
FACT0057
FACT0058
FACT0059
FACT0060
FACT0061
FACT0062
FACT0063
FACT0064
FACT0065
FACT0066
FACT0067
FACT0068
FACT0069
FACT0070
FACT0071
FACT0072
```

```
C  LNZE IS A SPECIAL CASE - NO SUM REQUIRED                          FACT0073
      IF (MCOL .EQ. M) GO TO 7                                       FACT0074
C  START SUM FROM GREATEST OF MROW OR ROW 'MCOL' LNZ COL NOS         FACT0075
      INIT = M                                                       FACT0076
      N = ILNZM1+MCOL                                                FACT0077
      IF (INTGR(N) .GT. M) INIT = INTGR(N)                           FACT0078
C  NO SUM IF ROW 'MCOL' HAS LEADING ZERCS UP TO THE DIAGONAL         FACT0079
      IF (INIT .EQ. MCOL) GO TO 7                                    FACT0080
C  ACCUMULATE THE SUM                                                FACT0081
      LAST = MCOL-1                                                  FACT0082
      DO 6 J = INIT,LAST                                             FACT0083
      KADRM = MM+J                                                   FACT0084
      KADRN = NN+J                                                   FACT0085
      JJ = IKOUM1+J                                                  FACT0086
      JJ = INTGR(JJ)+J                                               FACT0087
    6 SUM = SUM+REAL(JJ)*REAL(KADRM)*REAL(KADRN)                     FACT0088
C  BRANCH TO SPECIAL ALGORITHM FOR DIAGONAL ENTRIES                  FACT0089
    7 IF (MCOL .EQ. MROW) GO TO 8                                    FACT0090
C  FOR OFF-DIAGONAL ENTRIES:                                         FACT0091
      KADR = MM+MCOL                                                 FACT0092
      NN = NN+MCOL                                                   FACT0093
      REAL(KADR) = (REAL(KADR)-SUM)/REAL(NN)                         FACT0094
      GO TO 12                                                       FACT0095
C  DIAGONAL ENTRY - TEST FOR SINGULARITY AND SEMI-DEFINITENESS       FACT0096
    8 MM = MM+MROW                                                   FACT0097
      IF (REAL(MM)-SUM .NE. 0.) GO TO 9                              FACT0098
      JROW = MROW                                                    FACT0099
      GO TO 14                                                       FACT0100
    9 IF (REAL(MM)-SUM .GT. 0.) GO TO 10                             FACT0101
      WRITE (KW,902) MROW                                            FACT0102
      MPD = 1                                                        FACT0103
C  CALCULATE ROUNDING ERROR                                          FACT0104
   10 TESTR = ABS((REAL(MM)-SUM)/REAL(MM))                           FACT0105
      IF (TESTR .GE. TEST) GO TO 11                                  FACT0106
      TEST = TESTR                                                   FACT0107
      IROW = MROW                                                    FACT0108
```

```
C     EVALUATE DIAGONAL ENTRY                                          FACTO109
11    REAL(MM) = REAL(MM)-SUM                                          FACTO110
12    CONTINUE                                                         FACTO111
13    CONTINUE                                                         FACTO112
C     SEMI-DEFINITENESS CHECKS AND ROUNDING ERROR OUTPUT               FACTO113
      IF (MPD .EQ. 0) WRITE (KW,903)                                   FACTO114
      IF (MPD .EQ. 1 .AND. NPD .EQ. 1) WRITE (KW,905)                  FACTO115
      IF (MPD .EQ. 1 .AND. NPD .EQ. 0) WRITE (KW,904)                  FACTO116
      IERR = -1.00001*ALOG10(TEST)                                     FACTO117
      WRITE (KW,907) IROW,IERR                                         FACTO118
      IF (IERR .GT. 5) WRITE (KW,908)                                  FACTO119
      IF (MPD .EQ. 1 .AND. NPD .EQ. 1) STOP                            FACTO120
      IF (IERR .GT. 5) STOP                                            FACTO121
      RETURN                                                           FACTO122
C     SINGULAR MATRIX                                                  FACTO123
14    WRITE (KW,906) JROW                                              FACTO124
      STOP                                                             FACTO125
      END                                                              FACTO126
```

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      SUBROUTINE ORK(LENGTH)                                         ORK  0001
C*****************************************************************ORK  0002
C FINITE ELEMENT ANALYSIS BASIC LIBRARY SUBROUTINE               ORK  0003
C AEROELASTIC AND STRUCTURES RESEARCH LABORATORY                 ORK  0004
C MASSACHUSETTS INSTITUTE OF TECHNOLOGY                          ORK  0005
C*****************************************************************ORK  0006
C                                                                ORK  0007
C     THIS SUBROUTINE CREATES                                    ORK  0008
C           1) THE LNZ VECTOR WHICH HOLDS THE COLUMN NUMBER      ORK  0009
C              OF THE LEADING NON-ZERO ENTRY IN EACH ROW OF      ORK  0010
C              THE ASSEMBLED "K" MATRIX                          ORK  0011
C           2) THE ADDRESS COUNT VECTOR WHICH HOLDS THE ABSOLUTE ORK  0012
C              ADDRESS OF THE DIAGONAL ENTRY FOR EACH ROW OF THE ORK  0013
C              ASSEMBLED K MATRIX, MINUS THE ROW NUMBER, AND     ORK  0014
C              CHECKS THAT K FITS IN THE /DATA/ VECTOR           ORK  0015
C                                                                ORK  0016
      DIMENSION REAL(10000),INTGR(10000)                         ORK  0017
      COMMON /IO/ KR, KW, KP, KT1, KT2, KT3                      ORK  0018
      COMMON /SIZE/ NET, NDT                                     ORK  0019
      COMMON /BEGIN/ ICCN,IKOUNT,ILNZ,IMASTR,IQ,IK              ORK  0020
      COMMON /END/ LCON,LKOUNT,LLNZ,LMASTR,LQ,LK                CRK  0021
      COMMON /DATA/ REAL                                         ORK  0022
C VERSION 1 RELEASE 1 AUGUST 1972                                ORK  0023
C                                                                ORK  0024
      EQUIVALENCE (REAL(1),INTGR(1))                             ORK  0325
C PRINT CONTROL                                                  ORK  0026
  100 FORMAT(49H0THE LENGTH OF THE "DATA" VECTOR FOR THIS CASE IS,I10,14HORK  0027
     1H WHICH EXCEEDS,I10,49H =THE MAXIMUM ALLOWED IN THE DIMENSION STATORK  0028
     2EMENT./39H EXECUTION TERMINATED IN SUBROUTINE ORK)          ORK  0029
  200 FORMAT(5X,3HROW,2X,13HLNZE CCL. NO.,2X,18HABS. ADR. OF DIAG.) ORK  0030
  300 FORMAT(I9,I10,I12)                                          ORK  0031
  400 FORMAT(10H0THERE ARE,I10,68H NON-ZERO ENTRIES IN "K". IF "K" WERE ORK  0032
     1FULLY POPULATED THERE WOULD BE,I10, 9H ENTRIES.,/,20X,15HTHE DENSIORK  0033
     2TY IS ,E15.6)                                               ORK  0034
      ILNZM1=ILNZ-1                                               ORK  0035
      IMSTM1=IMASTR-1                                             ORK  0036
      IKOUM1=IKOUNT-1
      NETM1=NET-1
```

```
      IMSTP1=IMASTR+1                                                    ORK 0037
C SET EACH LNZ COLUMN NC = ROW NO (DIAGONAL MATRIX)                      ORK 0038
      DO 30 IROW=1,NDT                                                   ORK 0039
      MSUB=ILNZM1+IROW                                                   ORK 0040
30    INTGK(MSUB)=IROW                                                   ORK 0041
C EXAMINE MASTER ASSEMBLY LIST, ONE ELEMENT AT A TIME, TO CREATE         ORK 0042
C THE LNZ VECTOR                                                         CRK 0043
      DO 20 LNUM=1,NET                                                   ORK 0044
      MADDR = IMSTM1+LNUM                                                ORK 0045
      MADDR = INTGK(MADDR)-1                                             ORK 0046
C CALCULATE NO. DOF IN THE ELEMENT BY DIFFERENCING POINTERS, OR ...      ORK 0047
      I = IMASTR+LNUM                                                    ORK 0048
      IF(LNUM.EQ.NET) GO TO 3                                            ORK 0049
      NDE = INTGR(I)-INTGR(I-1)                                          ORK 0050
      GO TO 4                                                            ORK 0051
C ... BEGIN BY ASSUMING THE LIST IS FILLED, FOR LAST ELEMENT             ORK 0052
3     NDE = LMASTR-INTGR(I-1)+1                                          ORK 0053
C INITIALIZE SMALLEST DOF NO. AT LARGEST POSSIBLE VALUE                  ORK 0054
4     ISMALL=NDT                                                         ORK 0055
C FIND SMALLEST MASTR NUMBER FOR THIS ELEMENT                            ORK 0056
      DO 5 JDOF=1,NDE                                                    ORK 0057
      INDEX=MADDR+JDOF                                                   ORK 0058
      IF(INTGR(INDEX).GT.ISMALL) GO TO 5                                 ORK 0059
C DISCONTINUE SEARCH IF A ZERO IS FOUND, INDICATING EXCESS STORAGE       ORK 0060
C AND PREMATURE END OF LIST FOR LAST ELEMENT                             ORK 0061
      IF(INTGR(INDEX).EQ.0) GO TO 6                                      ORK 0062
      ISMALL=INTGR(INDEX)                                                ORK 0063
5     CONTINUE                                                           ORK 0064
C FIND COLUMN NUMBER OF LEADING NON-ZERO ENTRY IN ROW                    ORK 0065
6     DO 10 JDOF=1,NDE                                                   ORK 0066
      INDEX=MADDR+JDOF                                                   ORK 0067
      INDEX=ILNZM1+INTGR(INDEX)                                          ORK 0068
C CHANGE LNZ COL NO ONLY IF NEW ONE IS LESS THAN OLD ONE                 ORK 0069
      IF(INTGR(INDEX).LT.ISMALL) GO TO 8                                 ORK 0070
      INTGR(INDEX)=ISMALL                                                ORK 0071
      GO TO 10                                                           ORK 0072
```

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C DISCONTINUE OPERATION IF EXCESS STORAGE IS DISCOVERED         ORK 0073
8     IF(INTGR(INDEX).EQ.0) GO TO 20                            ORK 0074
10    CONTINUE                                                  ORK 0075
20    CONTINUE                                                  ORK 0076
C CREATE ADDRESS COUNT VECTOR                                   ORK 0077
      INTGR(LKOUNT) = LK                                        ORK 0078
      INDEX=LKOUNT                                              ORK 0079
      DO 40 IROW=2,NDT                                          ORK 0080
      I = ILN2M1+IROW                                           ORK 0081
      INTGR(INDEX+1) = INTGR(INDEX)+IROW+1-INTGR(I)             ORK 0082
40    INDEX=INDEX+1                                             ORK 0083
C ADDRESS COUNT VECTOR NOW CONTAINS ABSOLUTE ADDRESS ONLY FOR THE  ORK 0084
C DIAGONAL ENTRIES, AND THUS INTGR(LKOUNT) = LK EXACTLY         ORK 0085
      IF (INTGR(LKOUNT) .LE. LENGTH) GO TO 50                   ORK 0086
      WRITE (KW,100) INTGR(LKOUNT), LENGTH                      ORK 0087
      STOP                                                      ORK 0088
50    LK = INTGR(LKOUNT)                                        ORK 0089
      WRITE(KW,200)                                             ORK 0090
      DO 60 IROW=1,NDT                                          ORK 0091
      I = ILN2M1+IROW                                           ORK 0092
      J = IKOUM1+IROW                                           ORK 0093
      WRITE (KW,300) IROW, INTGR(I), INTGR(J)                   ORK 0094
C REPLACE THE ABS. ADDRESS OF DIAG. BY (ABS. ADDRESS - ROW NO.) ORK 0095
      INTGR(J) = INTGR(J)-IROW                                  ORK 0096
60    CONTINUE                                                  ORK 0097
      NENTRY = INTGR(LKOUNT)+NDT-LK+1                            ORK 0098
      INDEX = (NDT*(NDT+1))/2                                   ORK 0099
      DENS = FLOAT(NENTRY)/FLOAT(INDEX)                         ORK 0100
      WRITE (KW,400) NENTRY, INDEX, DENS                        ORK 0101
C ZERO THE FORCE/DISPLACEMENT VECTOR AND THE K MATRIX BLOCK     ORK 0102
      DO 70 I=10,LK                                             ORK 0103
70    REAL(I)=0.                                                ORK 0104
      RETURN                                                    ORK 0105
      END                                                       ORK 0106
```

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      SUBROUTINE ROTATE(NODE,IROW,JROW,KROW,ZANGLE,YANGLE,XANGLE)       ROTA0001
C*****************************************************************************ROTA0002
C FINITE ELEMENT ANALYSIS BASIC LIBRARY SUBROUTINE                      ROTA0003
C AEROELASTIC AND STRUCTURES RESEARCH LABORATORY                        ROTA0004
C MASSACHUSETTS INSTITUTE OF TECHNOLOGY                                 ROTA0005
C*****************************************************************************ROTA0006
      DIMENSION REAL(10000),INTGR(10000)                               ROTA0007
      DIMENSION A(3,3), B(3,3), IR(3), INFC(10), KDR(3), PERM(3,3),     ROTA0008
     2 PROTOR(3,3), ROTOR(3,3), TEMPC(3,3), TEMPR(3), TEMPC(3), TEMPRR(3)ROTA0009
      COMMON /IU/ KK, KW, KP, KT1, KT2, KT3                             ROTA0010
      COMMON /SIZE/ NET, NDT                                            ROTA0011
      COMMON /BEGIN/ ICON,IKOUNT,ILNZ,IMASTR,IQ,IK                      ROTA0012
      COMMON /DATA/ REAL                                                ROTA0013
C VERSION 1 RELEASE 1 AUGUST 1972                                       ROTA0014
C                                                                       ROTA0015
C                                                                       ROTA0016
      EQUIVALENCE (REAL(1), INTGR(1))                                   ROTA0017
C PRINT CONTROL                                                         ROTA0018
  901 FORMAT(1H0,41X,26HROTATION REQUESTED AT NODE,I10,/,24X,29HASSOCIATROTA0019
     1ED MASTER NUMBERS FOR,8X,5HFIRST,6X,6HSECOND,7X,5HTHIRD,3X,3HDOF,/ROTA0020
     2,54X,3(2X,I10),/,27X,27HEULER ANGLES IN DEGREES ARE,8X,7HTHETA Z,5ROTA0021
     3X,7HTHETA Y,5X,7HTHETA X,/,57X,3(2X,E10.3))                       ROTA0022
  902 FORMAT (1X,79HHALF TRANSFORM (K)*(THETA TRANSPOSE) WAS APPLIED TO ROTA0023
     1THE FOLLOWING DOF BELOW ROW,I10,16H IN THE K MATRIX)              ROTA0024
  903 FORMAT (10(2X,I10))                                              ROTA0025
  904 FORMAT(1X,103HUABOVE REQUEST INCONSISTENT DUE TO INEQUALITY OF LEADINROTA0026
     1G NON-ZERO ENTRY COLUMN NUMBERS OF ARGUMENT ROWS,/,9X,3HROW,1X,11HROTA0027
     2LNZE COL NO,/,3(2X,I10,2X,I10,/))                                 ROTA0028
  905 FORMAT (1X,41HEXECUTION TERMINATED IN SUBROUTINE ROTATE)          ROTA0029
  906 FORMAT (54HOABOVE REQUEST INCONSISTENT DUE TO REPEATED ROW NUMBER)ROTA0030
  907 FORMAT(73HOABOVE REQUEST INVALID DUE TO INCORRECT LOCATION OF UNDEROTA0031
     1FINED ROW NUMBERS)                                               ROTA0032
      IR(1)=IROW                                                        ROTA0033
      IR(2)=JROW                                                        ROTA0034
      IR(3)=KROW                                                        ROTA0035
C PRINT ENTRY MESSAGE                                                   ROTA0036
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      NNODE=IABS(NODE)
      WRITE (KW,901) NNODE, (IR(I), I = 1,3), ZANGLE, YANGLE, XANGLE       ROTA0037
C ESTABLISH OUF ROTATION LIMIT                                            ROTA0038
      LIMIT = 3                                                           ROTA0039
      IF (KROW .LE. 0) LIMIT = 2                                          ROTA0040
      IF(IROW .GT. 0 .AND. JROW .GT. 0) GC TO 100                         ROTA0041
      WRITE(KW,907)                                                       ROTA0042
      WRITE(KW,905)                                                       ROTA0043
      STOP                                                                ROTA0044
C GET LNZE COL NOS AND TEST FOR EQUALITY                                  ROTA0045
100   DO 1 I = 1,LIMIT                                                    ROTA0046
      J = ILN2+IR(I)-1                                                    ROTA0047
    1 INFO(I) = INTGR(J)                                                  ROTA0048
      IF (LIMIT .EQ. 3 .AND. INFO(1) .EQ. INFO(2) .AND. INFO(2) .EQ.      ROTA0049
    2 INFO(3)) GO TO 2                                                    ROTA0050
      IF (LIMIT .EU. 2 .AND. INFO(1) .EQ. INFO(2)) GO TO 2               ROTA0051
C SOMETHING WRONG - PRINT ERROR MESSAGE                                   ROTA0052
      WRITE (KW,904) (IR(I), INFO(I), I = 1,LIMIT)                        ROTA0053
      WRITE (KW,905)                                                      ROTA0054
      STOP                                                                ROTA0055
C CHECK FUR KEYPUNCH ERROR IN ARGUMENT LIST                              ROTA0056
    2 IF (IROW .NE. JROW .AND. JROW .NE. KROW .AND. KROW .NE. IROW)      ROTA0057
    2 GO TO 3                                                             ROTA0058
      WRITE (KW,906)                                                      ROTA0059
      WRITE (KW,905)                                                      ROTA0060
      STOP                                                                ROTA0061
C OK TO PROCEED - FORM ROTOR MATRIX AFTER CONVERTING                     ROTA0062
C ANGLES TO RADIAN MEASURE                                                ROTA0063
    3 Z = 3.14159*ZANGLE/180.                                            ROTA0064
      Y = 3.14159*YANGLE/180.                                            ROTA0065
      X = 3.14159*XANGLE/180.                                            ROTA0066
      ROTOR(1,1) = COS(Y)*COS(Z)                                          ROTA0067
      ROTOR(1,2) = COS(Y)*SIN(Z)                                          ROTA0068
      ROTOR(1,3) = SIN(Y)                                                 ROTA0069
      ROTOR(2,2) = COS(X)*COS(Z)-SIN(X)*SIN(Y)*SIN(Z)                     ROTA0070
      ROTOR(2,3) = SIN(X)*COS(Y)                                          ROTA0071
                                                                          ROTA0072
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      ROTOR(3,3) = COS(X)*COS(Y)                                          ROTA0073
      ROTOR(2,1) = -ROTOR(1,2)                                            ROTA0074
      ROTOR(3,1) = -ROTOR(1,3)                                            ROTA0075
      ROTOR(3,2) = -ROTOR(2,3)                                            ROTA0076
C     SKIP K MATRIX SECTION IF A RE-ROTATION                             ROTA0077
      IF(NODE .LT. 0) GO TO 38                                            ROTA0078
C     ESTABLISH LEAST AND LARGEST DOF NUMBERS FROM ARGUMENT LIST         ROTA0079
      LEAST = IROW                                                        ROTA0080
      LARGST = JROW                                                       ROTA0081
      DO 4 I = 1,LIMIT                                                    ROTA0082
      IF (IR(I) .LT. LEAST) LEAST = IR(I)                                ROTA0083
      IF (IR(I) .GT. LARGST) LARGST = IR(I)                              ROTA0084
    4 CONTINUE                                                           ROTA0085
C     APPLY (ROTOR)*(K) TO ROWS IROW, JROW, KROW - ONE COLUMN AT A TIME  ROTA0086
C     FROM LN2E TO LEAST-1 -- SKIP SECTION IF LEAST ROW = 1ST ROW        ROTA0087
      INIT = INFO(1)                                                      ROTA0088
      IF (LEAST .EQ. 1 .OR. INIT .EQ. LEAST) GO TO 8                     ROTA0089
      LAST = LEAST-1                                                      ROTA0090
      DO 7 MCOL = INIT,LAST                                               ROTA0091
C     GET ENTRIES OF K MATRIX INTO TEMPR, ZERO TEMPC                     ROTA0092
      DO 5 I = 1,LIMIT                                                    ROTA0093
      J = IKOUNT+IR(I)-1                                                  ROTA0094
      KADR = INTGR(J)+MCOL                                                ROTA0095
      TEMPR(I) = REAL(KADR)                                               ROTA0096
    5 TEMPC(I) = 0.                                                       ROTA0097
C     APPLY TRANSFORM                                                     ROTA0098
      DO 6 I = 1,LIMIT                                                    ROTA0099
      DO 6 J = 1,LIMIT                                                    ROTA0100
    6 TEMPC(I) = TEMPC(I)+ROTOR(I,J)*TEMPR(J)                            ROTA0101
C     RESTORE TRANSFORMED COLUMN TO K MATRIX                             ROTA0102
      DO 7 I = 1,LIMIT                                                    ROTA0103
      J = IKOUNT+IR(I)-1                                                  ROTA0104
      KADR = INTGR(J)+MCOL                                                ROTA0105
    7 REAL(KADR) = TEMPC(I)                                               ROTA0106
C     TEST FOR ROWS INTERVENING BETWEEN IROW, JROW AND KROW             ROTA0107
    8 IF (LARGST .EQ. LEAST+LIMIT-1) GO TO 26                            ROTA0108
```

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C*************************************************************************ROTA0109
C SPECIAL ALGORITHMS FOR INTERVENING ROWS                                ROTA0110
C*************************************************************************ROTA0111
C PRINT HEADING                                                          ROTA0112
      WRITE (KW,902) LEAST                                               ROTA0113
C FIND VALUE OF MIDDLE ARGUMENT ROW IF 3 DOF ARE BEING ROTATED           ROTA0114
      IF (LIMIT .EQ. 2) GO TO 10                                         ROTA0115
      DO 9 I = 1,3                                                       ROTA0116
      IF (IR(I) .NE. LEAST .AND. IR(I) .NE. LARGST) MIDDLE = IR(I)       ROTA0117
    9 CONTINUE                                                           ROTA0118
C TEST TO SEE IF ARGUMENT ROWS WERE GIVEN IN ASCENDING ORDER. IF         ROTA0119
C NOT, A PERMUTATION SIMILARITY TRANSFORM MUST BE APPLIED TO THE         ROTA0120
C ROTOR MATRIX TO TRANSFORM THE INTERVENING ROWS PROPERLY               ROTA0121
      IF (IROW .LT. JROW .AND. JROW .LT. KROW) GO TO 15                  ROTA0122
   10 IF (KROW .EQ. 0 .AND. IROW .LT. JROW) GO TO 15                    ROTA0123
C PREPARE STORAGE                                                        ROTA0124
      DO 11 I = 1,LIMIT                                                  ROTA0125
      DO 11 J = 1,LIMIT                                                  ROTA0126
      A(I,J) = 0.                                                        ROTA0127
      PROTOR(I,J) = 0.                                                   ROTA0128
   11 PERM(I,J) = 0.                                                     ROTA0129
      DO 12 I = 1,LIMIT                                                  ROTA0130
      IF (IR(I) .EQ. LEAST) PERM(1,I) = 1.                              ROTA0131
      IF (IR(I) .EQ. LARGST) PERM(LIMIT,I) = 1.                         ROTA0132
      IF (LIMIT .EQ. 2) GO TO 12                                         ROTA0133
      IF (IR(I) .EQ. MIDDLE) PERM(2,I) = 1.                             ROTA0134
   12 CONTINUE                                                           ROTA0135
      DO 13 I = 1,LIMIT                                                  ROTA0136
C APPLY SIMILARITY TRANSFORM                                             ROTA0137
      DO 13 J = 1,LIMIT                                                  ROTA0138
      DO 13 K = 1,LIMIT                                                  ROTA0139
   13 A(I,J) = A(I,J)+PERM(I,K)*ROTOR(K,J)                              ROTA0140
      DO 14 I = 1,LIMIT                                                  ROTA0141
      DO 14 J = 1,LIMIT                                                  ROTA0142
      DO 14 K = 1,LIMIT                                                  ROTA0143
   14 PROTOR(I,J) = PROTOR(I,J)+A(I,K)*PERM(J,K)                        ROTA0144
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      GO TO 17                                                          ROTA0145
C PERMUTATION NOT REQUIRED                                              ROTA0146
   15 DO 16 I = 1,3                                                     ROTA0147
      DO 16 J = 1,3                                                     ROTA0148
   16 PROTOR(I,J) = ROTOR(I,J)                                          ROTA0149
C INITIALIZE INFO VECTOR FILL INDEX AND ESTABLISH LOOP LIMITS FOR       ROTA0150
C INTERVENING ROWS                                                     ROTA0151
   17 INFOX = 0                                                         ROTA0152
      INIT = LEAST+1                                                    ROTA0153
C LAST ROW TO BE DONE DEPENDS ON WHETHER 2 OR 3 DOF ARE BEING ROTATED   ROTA0154
      IF (LIMIT .EQ. 2) LAST = LARGST-1                                 ROTA0155
      IF (LIMIT .EQ. 3) LAST = MIDDLE-1                                 ROTA0156
C LOOP OVER REMAINING COLUMNS, (PROTOR)*(K), AND FIRST ROWS,            ROTA0157
C (K)*(PROTOR-T), FROM INIT TO LAST -- SKIP IF MIDDLE DOF = LEAST+1     ROTA0158
C WHEN 3 DOF ARE BEING ROTATED                                         ROTA0159
      IF (LIMIT .EQ. 3 .AND. LAST .EQ. LEAST) GO TO 25                  ROTA0160
C THIS IS THE RE-ENTRY POINT IF 3 DOF ARE BEING ROTATED AND ROWS        ROTA0161
C INTERVENE BETWEEN MIDDLE AND LARGEST                                 ROTA0162
   18 DO 24 KROW = INIT,LAST                                            ROTA0163
C CHECK LINE COL NO OF CURRENT ROW -- IF .GT. LOWER ROTATED DOF,        ROTA0164
C WHICH = INIT-1, K MATRIX CONTAINS STORED OR NON-STORED ZEROS AND      ROTA0165
C BOTH ROW AND COLUMN TRANSFORM CAN BE SKIPPED                         ROTA0166
      J = ILN2+MKROW-1                                                  ROTA0167
      IF (INTGR(J) .GT. INIT-1) GO TO 24                                ROTA0168
C INCREMENT INFO FILL INDEX AND STORE CURRENT ROW NUMBER                ROTA0169
      INFOX = INFOX+1                                                   ROTA0170
      INFO(INFOX) = MROW                                                ROTA0171
C IF TEN ROW NUMBERS HAVE BEEN COLLECTED, PRINT THEM AND RESET INDEX    ROTA0172
      IF (INFOX .LT. 10) GO TO 19                                       ROTA0173
      WRITE (KW,903) (INFO(J), J = 1,10)                                ROTA0174
      INFOX = 0                                                         ROTA0175
C GET CORRECT ABSOLUTE ADDRESSES FOR K MATRIX ENTRIES, REPLACING COLUMN ROTA0176
C SECTORS ABOVE LOWER TRIANGLE BY ROW SECTORS BELOW, AND VICE-VERSA     ROTA0177
   19 J = IKOUNT+MKROW-1                                                ROTA0178
      KDR(I) = INTGR(J)                                                 ROTA0179
      J = IKOUNT+LAST                                                   ROTA0180
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         KDR(2) = INTGR(J)+MROW                                          ROTA0181
         IF (INIT .EQ. LEAST+1) GO TO 20                                 ROTA0182
         KDR(2) = KDR(1)+INIT-1                                          ROTA0183
   20    KDR(1) = KDR(1)+LEAST                                          ROTA0184
         IF (LIMIT .EQ. 2) GO TO 21                                      ROTA0185
         J = IKOUNT+LARGST-1                                            ROTA0186
         KDR(3) = INTGR(J)+MROW                                          ROTA0187
C PICK UP K MATRIX ENTRIES IN TEMPC, TEMPR AND ZERO TEMPCC, TEMPRR       ROTA0188
   21    DO 22 I = 1,LIMIT                                               ROTA0189
         KADR = KDR(I)                                                   ROTA0190
         TEMPC(I) = REAL(KADR)                                           ROTA0191
         TEMPR(I) = REAL(KADR)                                           ROTA0192
         TEMPCC(I) = 0.                                                  ROTA0193
   22    TEMPRR(I) = 0.                                                  ROTA0194
C APPLY TRANSFORMS (PROTOR)*(TEMPC) AND (TEMPR)*(PROTOR-TRANSPOSE)       ROTA0195
         DO 23 I = 1,LIMIT                                               ROTA0196
         DO 23 J = 1,LIMIT                                               ROTA0197
         TEMPCC(I) = TEMPCC(I)+PROTOR(I,J)*TEMPC(J)                      ROTA0198
   23    TEMPRR(I) = TEMPRR(I)+TEMPR(J)*PROTOR(I,J)                      ROTA0199
C RESTORE TRANSFORMED VALUES TO K MATRIX                                 ROTA0200
         KADR = KDR(1)                                                   ROTA0201
         REAL(KADR) = TEMPCC(1)                                          ROTA0202
         KADR = KDR(2)                                                   ROTA0203
         IF (INIT .EQ. LEAST+1) REAL(KADR) = TEMPRR(2)                   ROTA0204
         IF (INIT .NE. LEAST+1) REAL(KADR) = TEMPCC(2)                   ROTA0205
         IF (LIMIT .EQ. 2) GO TO 24                                      ROTA0206
         KADR = KDR(3)                                                   ROTA0207
         REAL(KADR) = TEMPRR(3)                                          ROTA0208
   24    CONTINUE                                                        ROTA0209
C PRINT PARTIALLY FILLED INFO VECTOR IF THERE IS AT LEAST ONE ROW.       ROTA0210
C RESET INDEX                                                            ROTA0211
         IF (INFOX .EQ. 0) GO TO 25                                      ROTA0212
         WRITE (KW,903) (INFO(I), I = 1,INFOX)                           ROTA0213
         INFOX = 0                                                       ROTA0214
C COMPLETION TEST                                                        ROTA0215
   25    IF (LIMIT .EQ. 2 .OR. INIT .NE. LEAST+1) GO TO 26               ROTA0216
```

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C CHECK FCR INTERVENING RUNS BETWEEN MIDDLE AND LARGEST          RCTA0217
      IF (MIDCLE+1 .EQ. LARGST) GO TO 26                          ROTA0218
C RESET LCCP LIMITS, PRINT NEW HEADING                           ROTA0219
      INIT = MIDCLE+1                                             ROTA0220
      LAST = LARGST-1                                             ROTA0221
      WRITE (K6,902) MIDDLE                                       RCTA0222
      GO TO 16                                                    ROTA0223
C******************************************************************ROTA0224
C END OF SPECIAL ALGORITHM SECTION                               ROTA0225
C******************************************************************ROTA0226
C FCR 2X2 CR 3X3 K MATRIX ENTRIES DIRECTLY ASSOCIATED WITH CCF BEING  ROTA0227
C RCTATEC, APPLY FULL SIMILARITY TRANSFCRM                       ROTA0228
   26 DO 28 I = 1,LIMIT                                           ROTA0229
      DO 28 J = 1,I                                               ROTA0230
      K = IKOUNT+IR(I)-1                                          RCTA0231
      KACR = INTGR(K)+IR(J)                                       ROTA0232
      IF (IR(I) .GT. IR(J)) GO TO 27                              RCTA0233
      K = IKOUNT+IR(J)-1                                          RCTA0234
      KACR = INTGK(K)+IR(I)                                       ROTA0235
C STCRE K MATRIX ENTRY IN A, ZERO B AND SYMMETRIZE               ROTA0236
   27 A(I,J) = REAL(KADK)                                         ROTA0237
      B(I,J) = 0.                                                 RCTA0238
      A(J,I) = A(I,J)                                             ROTA0239
   28 B(J,I) = 0.                                                 ROTA0240
C APPLY HALF CF TRANSFORM, (KUTUR)*(A)                           RCTA0241
      DC 29 I = 1,LIMIT                                           ROTA0242
      DC 29 J = 1,LIMIT                                           ROTA0243
      DC 29 K = 1,LIMIT                                           ROTA0244
   29 B(I,J) = B(I,J)+RUTUR(I,K)*A(K,J)                           ROTA0245
C FREPARE STORAGE                                                RCTA0246
      CC 30 I = 1,LIMIT                                           RCTA0247
      CC 30 J = 1,LIMIT                                           ROTA0248
   30 A(I,J) = 0.                                                 RCTA0249
C APPLY 2ND HALF CF TRANSFORM, (B)*(RCTOR-TRANSPCSE)             RCTA0250
      CG 31 I = 1,LIMIT                                           ROTA0251
      CC 31 J = 1,LIMIT                                           ROTA0252
```

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      DO 31 K = 1,LIMIT                                          ROTA0253
   31 A(I,J) = A(I,J)+B(I,K)*ROTOR(J,K)                          ROTA0254
C RESTORE TRANSFORMED VALUES TO K MATRIX                         ROTA0255
      DO 32 I = 1,LIMIT                                          ROTA0256
      DO 32 J = 1,I                                              ROTA0257
      K = IKOUNT+IR(I)-1                                         ROTA0258
      KADR = INTGR(K)+IR(J)                                      ROTA0259
      IF (IR(I) .GE. IR(J)) GO TO 32                             ROTA0260
      K = IKOUNT+IR(J)-1                                         ROTA0261
      KADR = INTGR(K)+IR(I)                                      ROTA0262
   32 REAL(KADR) = A(I,J)                                        ROTA0263
C APPLY (K)*(ROTOR-TRANSPOSE) TO ROWS BELOW LARGEST IF LNZE COL NOS  ROTA0264
C ARE .LE. LEAST. (IF LNZE COL NO IS .GT. LEAST, THE ROW CONTAINS    ROTA0265
C STORED OR NON-STORED ZEROS IN COLUMNS IROW, JROW, KROW AND CAN     ROTA0266
C BE SKIPPED.) -- SKIP THIS SECTION IF LARGEST DOF = NDT             ROTA0267
      IF (LARGST .EQ. NDT) GO TO 38                              ROTA0268
C PRINT NEW HEADING. RESET LOOP LOWER LIMIT AND INFO INDEX       ROTA0269
      WRITE (KW,902) LARGST                                      ROTA0270
      INIT = LARGST+1                                            ROTA0271
      INFOX = 0                                                  ROTA0272
      DO 37 MROW = INIT,NDT                                      ROTA0273
C CHECK LNZE COL NO                                              ROTA0274
      J = ILNZ+MROW-1                                            ROTA0275
      IF (INTGR(J) .GT. LEAST) GO TO 37                          ROTA0276
C INCREMENT INFO INDEX, STORE ROW NUMBER                         ROTA0277
      INFOX = INFOX+1                                            ROTA0278
      INFO(INFOX) = MROW                                         ROTA0279
C IF INFO VECTOR IS FILLED, PRINT AND RESET INDEX                ROTA0280
      IF (INFOX .LT. 10) GO TO 33                                ROTA0281
      WRITE (KW,903) (INFO(I), I = 1,10)                         ROTA0282
      INFOX = 0                                                  ROTA0283
C GET ENTRIES OF K MATRIX INTO TEMPC, ZERO TEMPR                 ROTA0284
   33 K = IKOUNT+MROW-1                                          ROTA0285
      K = INTGR(K)                                               ROTA0286
      DO 34 I = 1,LIMIT                                          ROTA0287
      KADR = K+IR(I)                                             ROTA0288
```

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```
          TEMPC(I) = REAL(KADR)
   34 TEMPR(I) = 0.
C APPLY TRANSFORM
          DO 35 I = 1,LIMIT
          DO 35 J = 1,LIMIT
   35 TEMPR(I) = TEMPR(I)+TEMPC(J)*ROTOR(I,J)
C RESTORE TO K
          DO 36 I = 1,LIMIT
          KADR = K+IR(I)
   36 REAL(KADR) = TEMPR(I)
   37 CONTINUE
C CHECK FOR PARTIALLY FILLED INFO VECTOR
          IF (INFOX .EQ. 0) GO TO 38
          WRITE (KM,903) (INFO(I), I = 1,INFOX)
C APPLY TRANSFORM (ROTOR)*(Q) TO ASSEMBLED ELEMENT EQUIVALENT
C NODAL FORCES
   38 DO 39 I = 1,LIMIT
          J = IQ+IR(I)-1
          TEMPR(I) = REAL(J)
   39 TEMPC(I) = 0.
          DO 40 I = 1,LIMIT
          DO 40 J = 1,LIMIT
   40 TEMPC(I) = TEMPC(I)+ROTOR(I,J)*TEMPR(J)
          DO 41 I = 1,LIMIT
          J = IQ+IR(I)-1
   41 REAL(J) = TEMPC(I)
          RETURN
          END
```

```
ROTA0289
ROTA0290
ROTA0291
ROTA0292
ROTA0293
ROTA0294
ROTA0295
ROTA0296
ROTA0297
ROTA0298
ROTA0299
ROTA0300
ROTA0301
ROTA0302
ROTA0303
ROTA0304
ROTA0305
ROTA0306
ROTA0307
ROTA0308
ROTA0309
ROTA0310
ROTA0311
ROTA0312
ROTA0313
ROTA0314
ROTA0315
ROTA0316
```

```
      SUBROUTINE SETUP(LENGTH,NCON,MASTRL)                             SETU0001
C*********************************************************************SETU0002
C FINITE ELEMENT ANALYSIS BASIC LIBRARY SUBROUTINE                    SETU0003
C AERUELASTIC AND STRUCTURES RESEARCH LABORATORY                      SETU0004
C MASSACHUSETTS INSTITUTE OF TECHNOLOGY                               SETU0005
C*********************************************************************SETU0006
      DIMENSION REAL(10000),INTGR(10000)                              SETU0007
      DIMENSION II(6), LI(6), KD(6)                                   SETU0008
      COMMON /IO/ KR, KW, KP, KT1, KT2, KT3                           SETU0009
      COMMON /SIZE/ NET, NDT                                          SETU0010
      COMMON /BEGIN/ ICON,IKOUNT,ILNZ,IMASTR,IQ,IK                    SETU0011
      COMMON /END/ LCON,LKOUNT,LLNZ,LMASTR,LQ,LK                      SETU0012
      COMMON /DATA/ REAL                                              SETU0013
C VERSION 1 RELEASE 1 AUGUST 1972                                     SETU0014
      EQUIVALENCE (REAL(1), INTGR(1))                                 SETU0015
      EQUIVALENCE (ICON,II(1)), (LCON,LI(1)), (KR,KD(1))              SETU0016
C PRINT CONTROL                                                       SETU0017
  901 FORMAT (1H0,53X,11HSETUP ENTRY,/,42X,35HUSER SPECIFICATIONS TO FEASETU0018
     1BL SYSTEM,/,1H0,51X,16HI/O DEVICE CODES,/,27X,6HREADER,5X,7HPRINTESETU0019
     2R,2X,10HCARD PUNCH,7X,5HTAPE1,7X,5HTAPE2,7X,5HTAPE3,/,21X,6(2X,11OSETU0020
     3),/,1H0,53X,12HPROBLEM SIZE,/,42X,25HTOTAL NUMBER OF ELEMENTS=,110SETU0021
     4,/,42X,25HTOTAL DEGREES OF FREEDOM=,110,/,1H0,51X,16HENTRY PARAMETSETU0022
     5ERS,/,39X,32HASSUMED LENGTH OF /DATA/ VECTOR=,110,/,32X,45HNUMBER SETU0023
     6OF DISPLACEMENT CONSTRAINS REQUESTED=,110,/,33X,44HNUMBER OF WORDSSETU0024
     7S REQUESTED FOR ASSEMBLY LIST=,110)                             SETU0025
  902 FORMAT (1H0,43X,31H/DATA/ VECTOR ADDRESS INDEX MAP,/,22X,11HCONSTRSETU0026
     1AINTS,2X,10HDG ABS ADR,1X,11HLNZE COL NO,1X,11HASMBLY LIST,2X,10HQSETU0027
     2/U VECTOR,4X,8HK MATRIX,/,16X,5HBEGIN,6(2X,110),/,18X,3HEND,5(2X,1SETU0028
     310),11X,1H+,/,1H0,36H+ LK IS CALCULATED BY SUBROUTINE ORK)      SETU0029
  903 FORMAT (40HOSTORAGE EXCEEDS LENGTH OF /DATA/ VECTOR,/,1X,34HL ENGTHSETU0030
     1 SUGGESTED FOR THIS PROBLEM=,112,6H WORDS,/,1X,40HEXECUTION TERMINSETU0031
     2ATED IN SUBROUTINE SETUP)                                      SETU0032
C*********************************************************************SETU0033
C REMOVE THIS FORMAT AND WRITE STATEMENT INDICATED BELOW IF FEABL     SETU0034
C HEADING IS NOT DESIRED                                              SETU0035
 1001 FORMAT(1H1 ,92(1HX)/),4(5H XXXX,84X,4HXXXX/),5H XXXX,7X,10(1HSETU0036
```

```
1F),5X,10(11E),9X,2HAA,9X,8(1HB),7X,3HLLL,14X,4HXXXX,7X,10(SETU0037
21HF),5X,10(1HE),8X,4HAAAA,8X,9(1HB),7X,3HLLL,6X,3HLLL,14X,4HXXXX/5H XXXX,7XSETU0038
3,3HFFF,12X,3HEEE,14X,6(1HA),7X,3HBBB,3X,4HBBBB,5X,3HLLL,14X,4HXXXXXSETU0039
4/5H XXXX,7X,3HFFF,12X,3HEEE,13X,3HAAA,6X,3HBBB,4X,3HBBB,5SETU0040
5X,3HLLL,14X,4HXXXX/5H XXXX,7X,3HFFF,12X,3HEEE,12X,3HAAA,4X,3HAAA,5SETU0041
6X,3HBBB,4X,3HBBB,5X,3HLLL,14X,4HXXXX/5H XXXX,7X,3HFFF,12X,3HEEE,12SETU0042
7X,3HAAA,4X,3HAAA,5X,3HBBB,3X,4HBBBB,5X,3HLLL,14X,4HXXXX/5H XXXX,7XSETU0043
8,8(1HF),7X,8(1HE),7X,3HAAA,4X,3HAAA,5X,9(1HB),6X,3HLLL,14X,4HXXXXSETU0044
95H XXXX,7X,8(1HF),7X,8(1HE),7X,10(1HA),5X,9(1HB),7X,10(1HL),SETU0045
AXX/5H XXXX,7X,3HFFF,12X,3HEEE,12X,1C(1HA),5X,3HBBB,3X,4HBBBB,5X,3HSETU0046
BLLL,14X,4HXXXX/5H XXXX,7X,3HFFF,12X,3HEEE,12X,3HAAA,4X,3HAAA,5X,3HSETU0047
CBBB,4X,3HBBB,5X,3HLLL,14X,4HXXXX/5H XXXX,7X,3HFFF,12X,3HEEE,12X,3HSETU0048
DAAA,4X,3HAAA,5X,3HBBB,3X,4HBBBB,5X,3HLLL,14X,4HXXXX/5H XXXX,7X,3HFSETU0049
EFF,12X,10(1HE),5X,3HAAA,4X,3HAAA,5X,9(1HB),7X,4HXXXX/5HSETU0050
F XXXX,7X,3HFFF,12X,10(1HE),5X,3HAAA,4X,3HAAA,5X,8(1HB),7X,10(1HL),SETU0051
G7X,4HXXXX/4(5H XXXX,84X,4HXXXX/),5H XXXX,24X,37HFINITE ELEMENT ANASETU0052
HLYSIS BASIC LIBRARY,23X,4HXXXX/4(5H XXXX,84X,4HXXXX/),3(1H ,92(1HXSETU0053
1)/),1H1)                                                      SETU0054
C ***************************************************************SETU0055
C PRINT ENTRY MESSAGE                                           SETU0056
C ***************************************************************SETU0057
C REMOVE THIS WRITE STATEMENT IF FEABL HEADING IS NOT DESIRED   SETU0058
      WRITE (KW,1001)                                           SETU0059
C ***************************************************************SETU0060
      WRITE (KW,901) (KD(I), I = 1,6), NET, NDT, LENGTH, NCON, MASTRL SETU0061
C CALCULATE ADDRESS INDEX VALUES                                SETU0062
      ICON = 1                                                  SETU0063
      LCON = NCON                                               SETU0064
      IKOUNT = LCON+1                                           SETU0065
      LKOUNT = LCON+NDT                                         SETU0066
      ILNZ = LKOUNT+1                                           SETU0067
      LLNZ = LKOUNT+NDT                                         SETU0068
      IMASTR = LLNZ+1                                           SETU0069
      LMASTR = LLNZ+MASTRL                                      SETU0070
      IQ = LMASTR+1                                             SETU0071
      LQ = LMASTR+NDT                                           SETU0072
```

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```
      IK = LU+1
C PRINT INDEX MAP                                              SETU0073
      WRITE (KW,902) (II(I), I = 1,6), (LI(I), I = 1,5)        SETU0074
C STORAGE BOUNDS TEST                                          SETU0075
      IF (LMASTK .LE. LENGTH) GO TO 1                          SETU0076
C STORAGE EXCEEDED - ESTIMATE REQUIRED LENGTH BASED ON LOWER TRIANGLE  SETU0077
C ESTIMATED POPULATION FACTOR                                  SETU0078
      TR = (NDT*(NDT+1))/2                                     SETU0079
      DENS = 0.5                                               SETU0080
      IF (NDT .GT. 200) DENS=0.3                               SETU0081
      IF (NDT .GT. 500) DENS=0.2                               SETU0082
      IF (NDT .GT. 1500) DENS = 0.15                           SETU0083
      IF (NDT .GT. 2000) DENS = 0.10                           SETU0084
      TR = TR*DENS                                             SETU0085
      LENGTH = LU+TR                                           SETU0086
      WRITE (KW,903) LENGTH                                    SETU0087
      STOP                                                     SETU0088
C ENOUGH STORAGE EXISTS TO GO THRU CRK                         SETU0089
    1 DO 2 I = ICON,LCON                                       SETU0090
    2 INTGR(I) = 0                                             SETU0091
      RETURN                                                   SETU0092
      END                                                      SETU0093
                                                               SETU0094
```

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      SUBROUTINE SIMULQ(ENERGY)                                         SIMU0001
C************************************************************************SIMU0002
C FINITE ELEMENT ANALYSIS BASIC LIBRARY SUBROUTINE                      SIMU0003
C AEROELASTIC AND STRUCTURES RESEARCH LABORATORY                        SIMU0004
C MASSACHUSETIS INSTITUTE OF TECHNOLOGY                                 SIMU0005
C************************************************************************SIMU0006
      DIMENSION REAL(10000), INTGR(10000)                               SIMU0007
      COMMON /IO/ KR, KW, KP, KT1, KT2, KT3                             SIMU0008
      COMMON /SIZE/ NET, NDT                                            SIMU0009
      COMMON /BEGIN/ ICCN, IKOUNT, ILNZ, IMASTR, IQ, IK                 SIMU0010
      COMMON /END/ LCON,LKOUNT,LLNZ,LMASTR,LQ,LK                        SIMU0011
      COMMON /DATA/ REAL                                                SIMU0012
C VERSION 1 RELEASE 1 AUGUST 1972                                       SIMU0013
C                                                                       SIMU0014
      EQUIVALENCE (REAL(1), INTGR(1))                                   SIMU0015
C PRINT CONTROL                                                         SIMU0016
  901 FORMAT (1H0,53X,18HSIMULQ ENTRY POINT,/,38H PRESCRIBED FORCE/DISPLSIMU0017
     2ACEMENT VECTOR: )                                                 SIMU0018
  902 FORMAT (30HODISPLACEMENT SOLUTION VECTOR: )                       SIMU0019
  903 FORMAT (1H0,2X,3HROW,10I12)                                       SIMU0020
  904 FORMAT (6H VALUE,10(2X,E10.3))                                    SIMU0021
  905 FORMAT (32HOSTRAIN ENERGY IN THE STRUCTURE=,E10.3)                SIMU0022
      IKOUM1 = IKOUNT-1                                                 SIMU0023
      ILNZM1 = ILNZ-1                                                   SIMU0024
      IQM1 = IQ-1                                                       SIMU0025
C PRINT ENTRY MESSAGE, SET PRINT SECTION CONTROL FLAG                   SIMU0026
      WRITE (KW,901)                                                    SIMU0027
      IFLAG = 0                                                         SIMU0028
C PRINT SECTION: DIVIDE Q VECTOR INTO GROUPS OF TEN                     SIMU0029
C PLUS A POSSIBLE REMAINDER                                             SIMU0030
      MOST = NDT/10                                                     SIMU0031
      LEFT = NDT-10*MOST                                                SIMU0032
      JBEG = 1                                                          SIMU0033
      NBEG = IQ                                                         SIMU0034
C RE-ENTRY POINT FOR SOLUTION OUTPUT                                    SIMU0035
    1 IF (NDT .LT. 10) GO TO 3                                          SIMU0036
```

```
      DO 2 I = 1,MUST
      JBEG = 1+IQ*(I-1)                                              SIMU0037
      JEND = JBEG+9                                                  SIMU0038
      NBEG = IQ+JBEG-1                                               SIMU0039
      NEND = NBEG+9                                                  SIMU0040
      WRITE (KW,903) (J, J = JBEG,JEND)                             SIMU0041
    2 WRITE (KW,904) (REAL(N), N = NBEG,NEND)                       SIMU0042
C CHECK FOR EXISTENCE OF REMAINDER                                  SIMU0043
      IF (LEFT .EQ. 0) GO TO 4                                       SIMU0044
      JBEG = JEND+1                                                  SIMU0045
      NBEG = NEND+1                                                  SIMU0046
    3 WRITE (KW,903) (J, J = JBEG,NDT)                              SIMU0047
      WRITE (KW,904) (REAL(N), N = NBEG,LC)                         SIMU0048
C CHECK CONTROL FLAG                                                SIMU0049
    4 IF (IFLAG .EQ. 0) GO TO 5                                      SIMU0050
      WRITE (KW,905) ENERGY                                          SIMU0051
      RETURN                                                         SIMU0052
C FORWARD SOLUTION - NO DIVISIONS, SO SKIP FIRST ROW ENTIRELY       SIMU0053
C INITIALIZE CONSTRAINT VECTOR AT 1ST NONZERO ENTRY                 SIMU0054
    5 DO 6 I = ICUN,LCCN                                             SIMU0055
      IF (INTGR(I) .EQ. 0) GO TO 6                                   SIMU0056
      GO TO 7                                                        SIMU0057
    6 CONTINUE                                                       SIMU0058
    7 NEXT = I                                                       SIMU0059
      IF (INTGR(NEXT) .EQ. 1) NEXT = NEXT+1                          SIMU0060
C SOLVE (A)(K) = (Q)                                                SIMU0061
      DO 10 MROW = 2,NDT                                             SIMU0062
C CHECK FOR CONSTRAINT TO SKIP ROW                                  SIMU0063
      IF (MROW .NE. INTGR(NEXT)) GO TO 8                             SIMU0064
C UPDATE NEXT                                                       SIMU0065
      NEXT = NEXT+1                                                  SIMU0066
      GO TO 10                                                       SIMU0067
C INITIALIZE SUM AND LOOP LIMITS FOR SUM                            SIMU0068
    8 SUM = 0.                                                       SIMU0069
      LAST = MROW-1                                                  SIMU0070
      M = ILN2+LAST                                                  SIMU0071
                                                                     SIMU0072
```

```
      M = INTGR(4)                                                    SIMU0073
      MM = IKOUNT+LAST                                                SIMU0074
      MM = INTGR(MM)                                                  SIMU0075
      DO 9 J = M,LAST                                                 SIMU0076
      KADR = MM+J                                                     SIMU0077
      N = IUM1+J                                                      SIMU0078
    9 SUM = SUM+REAL(KADR)*REAL(N)                                    SIMU0079
C SUBTRACT SUM FROM U                                                 SIMU0080
      N = IUM1+MROW                                                   SIMU0081
      REAL(N) = REAL(N)-SUM                                           SIMU0082
   10 CONTINUE                                                        SIMU0083
C SOLVE (D)(P) = (K) AND CALCULATE ENERGY                             SIMU0084
      N = IU-1                                                        SIMU0085
      ENERGY = 0.                                                     SIMU0086
      DO 11 MROW = 1,NDT                                              SIMU0087
      KADR = IKOUN1+MROW                                              SIMU0088
      KADR = INTGR(KADR)+MROW                                         SIMU0089
      N = N+1                                                         SIMU0090
      REAL(N) = REAL(N)/REAL(KADR)                                    SIMU0091
   11 ENERGY = ENERGY+REAL(KADR)*(REAL(N)**2)                         SIMU0092
      ENERGY = 0.5*ENERGY                                             SIMU0093
C BACK SOLUTION - NO DIVISIONS, SO SKIP LAST ROW ENTIRELY            SIMU0094
      NEXT = LCON                                                     SIMU0095
      IF (INTGR(NEXT) .EQ. NDT) NEXT = NEXT-1                         SIMU0096
C LOOP BACKWARDS OVER REMAINING ROWS                                  SIMU0097
      DO 14 I = 2,NDT                                                 SIMU0098
      MROW = NDT+1-I                                                  SIMU0099
C CHECK FOR CONSTRAINT TO SKIP ROW                                    SIMU0100
      IF (MROW .NE. INTGR(NEXT)) GO TO 12                             SIMU0101
C UPDATE NEXT                                                         SIMU0102
      NEXT = NEXT-1                                                   SIMU0103
      IF (NEXT .LT. ICON) NEXT = LCON                                 SIMU0104
      GO TO 14                                                        SIMU0105
C INITIALIZE SUM AND LOWER LOOP LIMIT                                 SIMU0106
   12 SUM = 0.                                                        SIMU0107
      INIT = MROW+1                                                   SIMU0108
```

```
      DO 13 J = INIT,NDT                                          SIMU0109
C CHECK IF LNZ COL NO OF ROW J EXCEEDS MROW - IF SO, SKIP         SIMU0110
      N = ILNZM1+J                                                SIMU0111
      IF (INTGR(N) .GT. MROW) GO TO 13                            SIMU0112
      KADR = IKUJM1+J                                             SIMU0113
      KADR = INTGR(KADR)+MROW                                     SIMU0114
      N = IJM1+J                                                  SIMU0115
      SUM = SUM+REAL(KADR)*REAL(N)                                SIMU0116
   13 CONTINUE                                                    SIMU0117
C SUBTRACT SUM FROM U                                             SIMU0118
      N = IJM1+MROW                                               SIMU0119
      REAL(N) = REAL(N)-SUM                                       SIMU0120
   14 CONTINUE                                                    SIMU0121
C PRINT SOLUTION HEADING, RESET CONTROL FLAG AND                 SIMU0122
C BRANCH TO OUTPUT SECTION                                        SIMU0123
      WRITE (KM,9002)                                            SIMU0124
      IFLAG = 1                                                   SIMU0125
      GO TO 1                                                     SIMU0126
      END                                                         SIMU0127
```